

Interpolation Search

Review talk, joint work with: S. Sioutas, C. Makris, C. Zaroliagis, T. Tsakalidis, K. Tsihlias

Results published at:

Conferences: ESA'03, ISAAC '05 & 09, ICALP'06, ICDT'10

Journals: JDA, Algorithmica

Talk :

Alexis C. Kaporis, Dept. Information & Communication Systems, U. Aegean, Karlovassi, Samos

Interpolation Search

BUT...

Interpolation Search

BUT... What is this "Interpolation" ?

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Easy things *should* come first:

Interpolation Search

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Easy things *should* come first: **Serial** search

Interpolation Search

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Suppose that Nature likes us...

Interpolation Search

BUT... What is this "Interpolation" ?

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File = [3, 5, 1, 0, 7, 9, 12, 16, 4]

target!

Interpolation Search

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File = [3, 5, 1, 0, 7, 9, 12, 16, 4]

target!

1 step!



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1 step! But...

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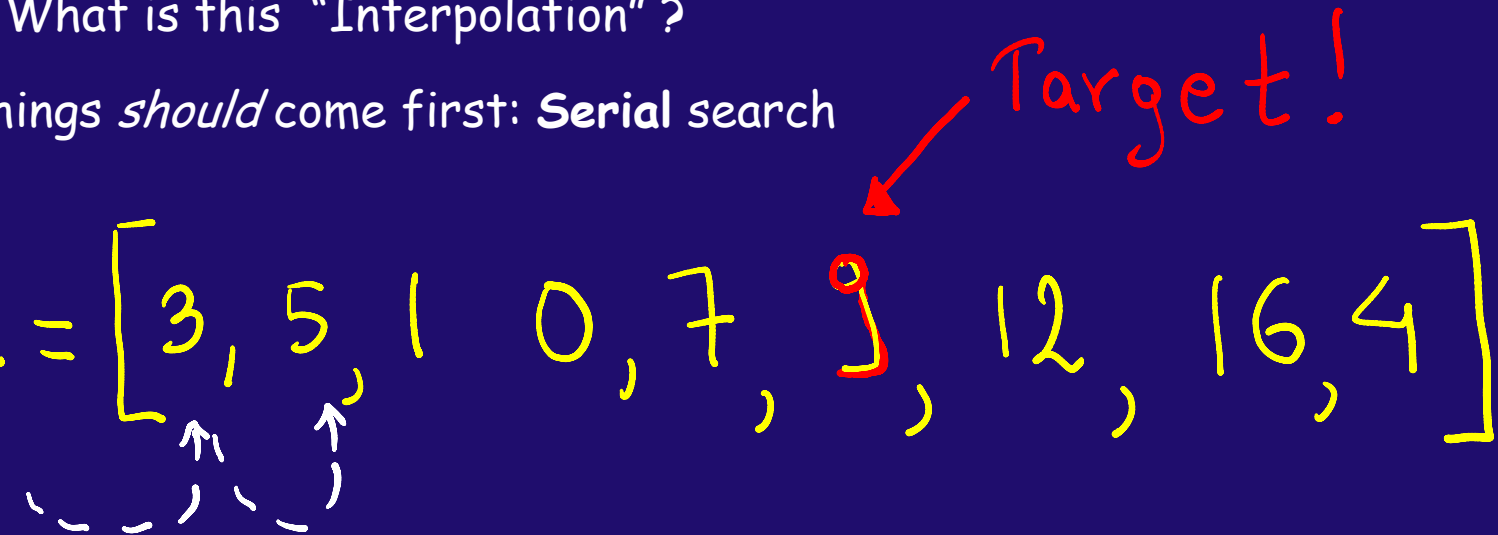
The image shows a handwritten array of numbers: [3, 5, 1, 0, 7, 9, 12, 16, 4]. A dashed white arrow starts from the first element '3' and points to the second element '5'. A red arrow points from the word 'Target!' to the element '9'.

Interpolation Search

BUT... What is this "Interpolation" ?

Easy things *should* come first: **Serial** search

File = [3, 5, 10, 7, 9, 12, 16, 4]



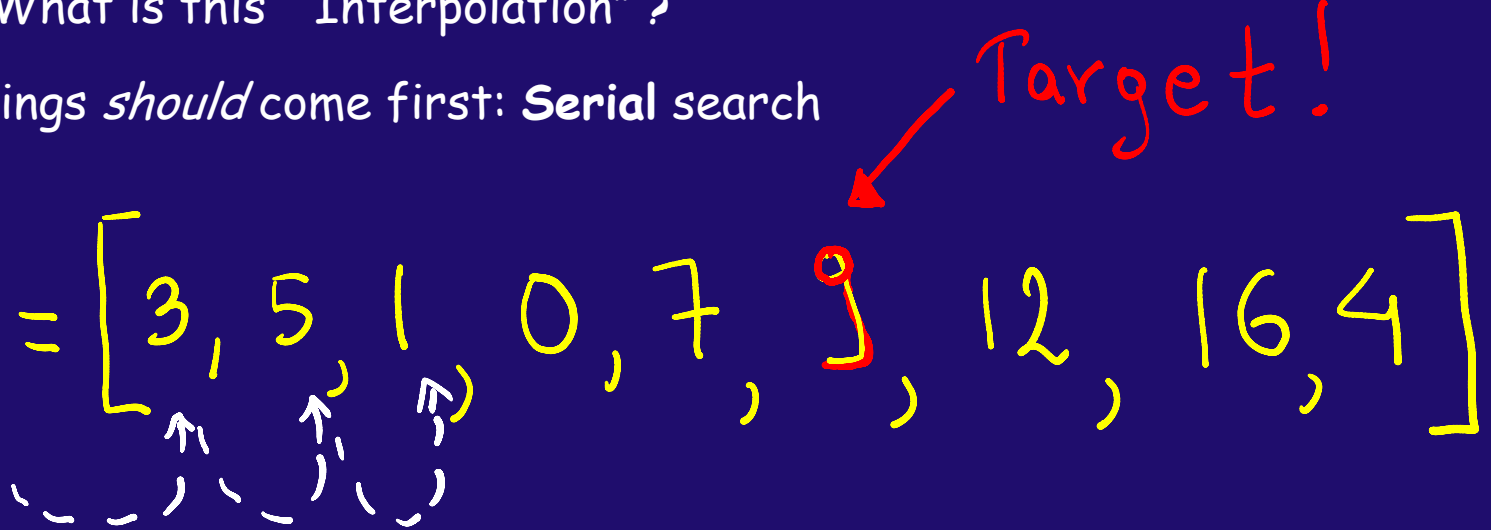
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Target!

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File = [3, 5, 1, 0, 7, 9, 12, 16, 4]

The diagram shows a sequence of numbers in brackets: 3, 5, 1, 0, 7, 9, 12, 16, 4. The number 9 is circled in red and has a red arrow pointing to it from the word "Target!" written in red above it. Below the first five numbers (3, 5, 1, 0, 7), there are dashed white arrows pointing upwards from a common baseline, representing a serial search process.

Interpolation Search

BUT... What is this "Interpolation" ?

Easy things *should* come first: **Serial** search

File = [3, 5, 1, 0, 7, 9, 12, 16, 4]

6 steps

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Deteriorates as $\Theta(n)$, as file size n increases

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Easy things *should* come first: **Serial** search

Suppose that she (Nature) still likes us

Interpolation Search

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Target!

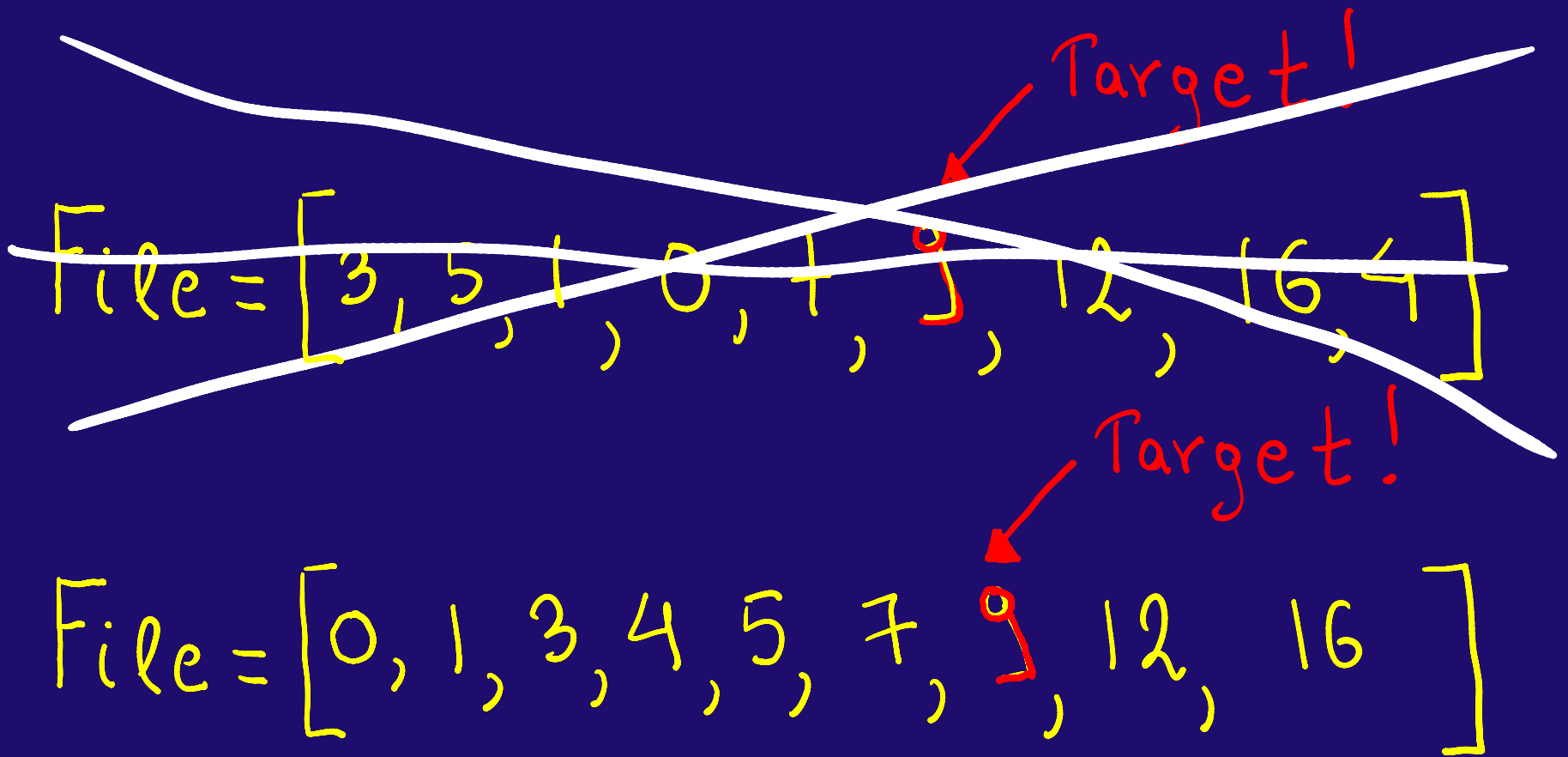
File = [0, 1, 3, 4, 5, 7, 9, 12, 16]

Target!

Interpolation Search

BUT... What is this "Interpolation" ?

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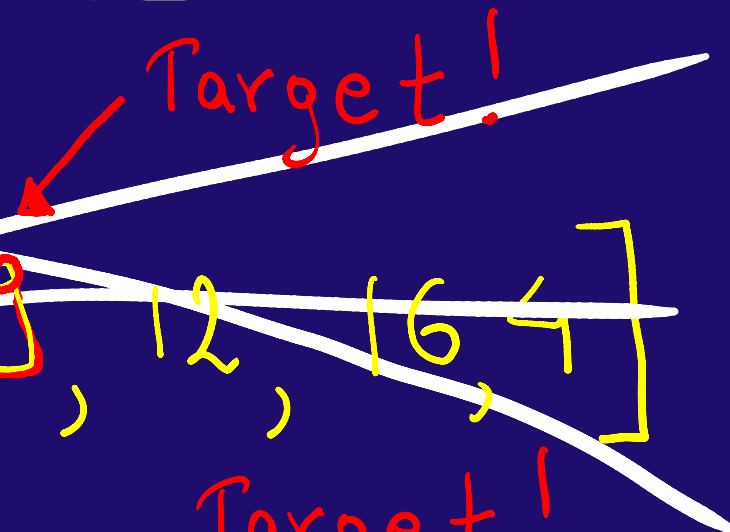
«Εν αρχη ην η Τάξηη», aka «assume ordered file»

Interpolation Search


BUT... What is this "Interpolation" ?

Easy things *should* come first: ~~Serial~~ search → Binary search

~~File = [3, 5, 1, 0, 7, 9, 12, 16, 4]~~



File = [0, 1, 3, 4, 5, 7, 9, 12, 16]




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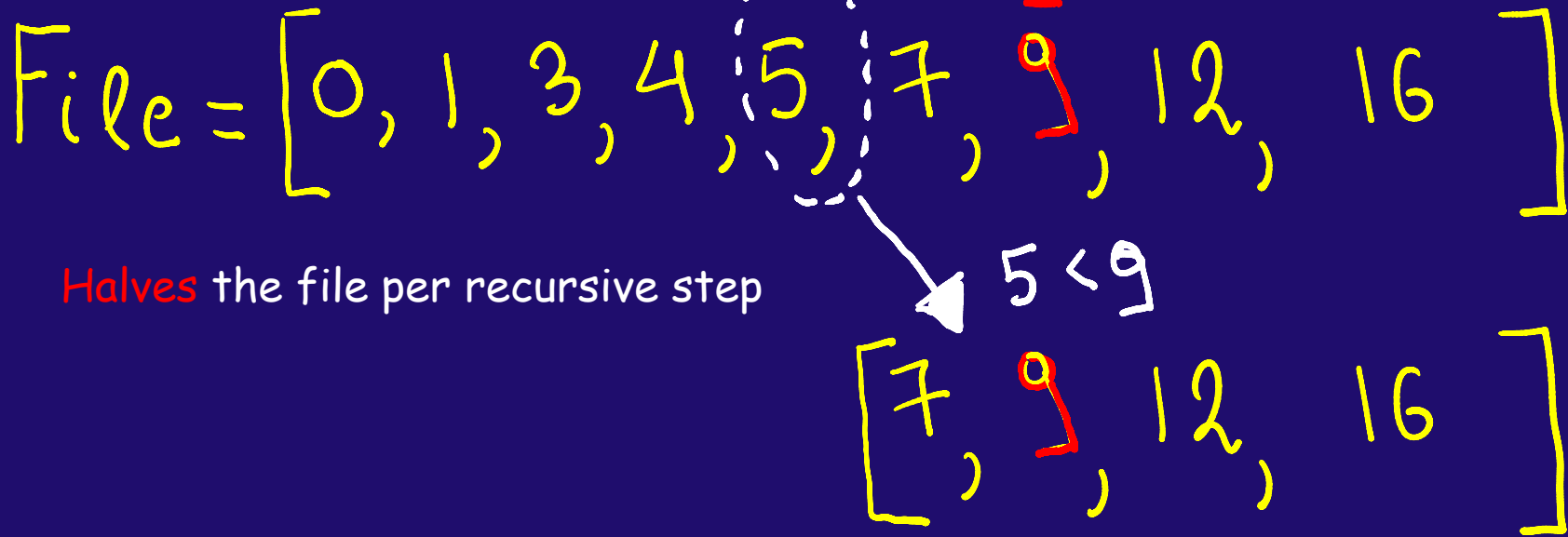


Halves the file per recursive step

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5 < 9
[7, 9]

Interpolation Search

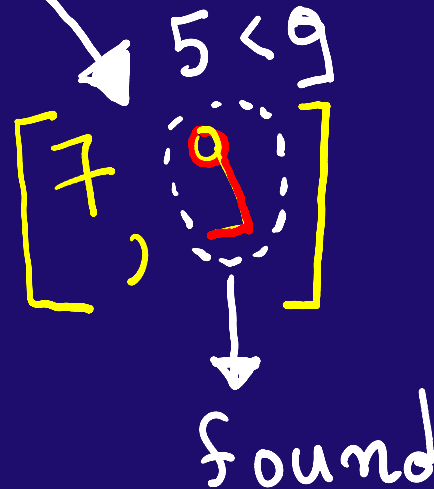
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File = [0, 1, 3, 4, 5, 7, 9, 12, 16]

Target!

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n → n/2 → n/2² → ... n/2ⁱ → ... 1

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$n \rightarrow n/2 \rightarrow n/2^2 \rightarrow \dots n/2^i \rightarrow \dots 1$

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Halves the file per recursive step

5 < 9
[7, 9]

$n \rightarrow n/2 \rightarrow n/2^2 \rightarrow \dots n/2^i \rightarrow \dots$ 1 $i = \log n$

Deteriorates as $\Theta(\log n)$, as file size n increases

Let us take a more **panoramic** view on...

...(random) inputs of Binary search

Let us take a more **panoramic** view on...

...(random) inputs of Binary search
...not a Πανάκεια wrt input distribution

Knuth's Example: a **lexicon** is a very nice and very ordered file

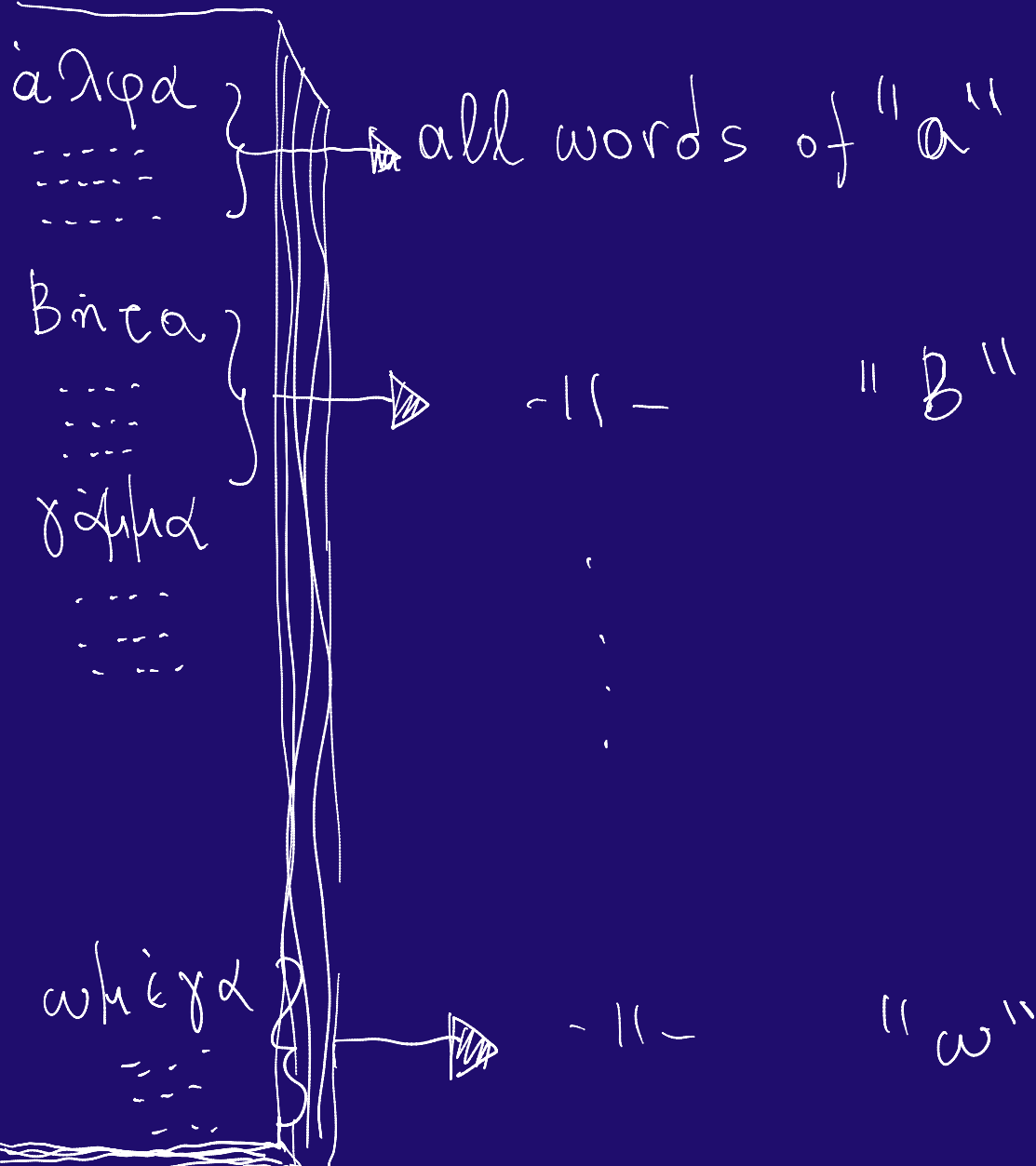
α λ φ δ

β η τ α

γ θ μ α

ω η ε γ δ

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γ ρ μ δ

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α λ φ α

β η τ α

γ α φ η α

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Search for "αύγουστος"

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Search for "αύγουστος"

Binary search will start around 12th letter = "μ"

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Search for "αύγουστος"

Binary search will start around 12th letter = "μ"

But "μ" is far away from "α"

distance

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α λ φ α

β η τ α

ζ α μ α

ω η ε γ α

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Search for "αύγουστος"

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But "ζ" is also away from "α"

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Binary search keeps walking to 3rd letter = "γ"

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And the story goes on ...

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α λ φ α

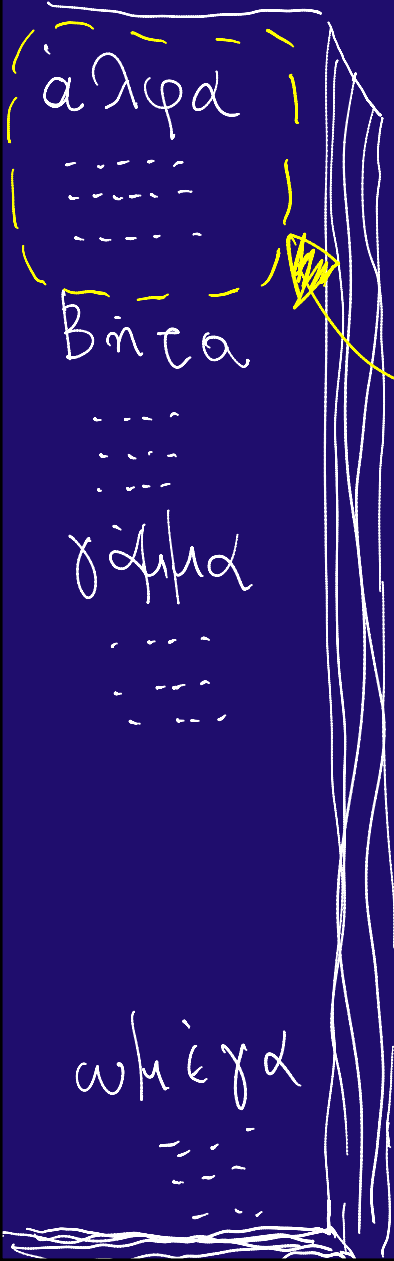
β η τ α

γ α φ η α

ω η ε γ α

How a secretariat (≠ Computer Scientist) would search for "αύγουστος"?

Knuth's Example: a **lexicon** is a very nice and very ordered file



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She expects "αύγουστος" to be near the front pages

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How a secretariat (≠ Computer Scientist) would search for "αύγουστος"?

She expects "αύγουστος" to be near the front pages

and opens the lexicon near the front pages, locating "αναψυχη" < "αύγουστος"

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β η τ α

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She observes "ν" is close to "ύ" and opens lexicon near to the current page, locating "αυγό"

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Now she is really close, within seconds locates "αύγουστος"

Knuth's Example: a **lexicon** is a very nice and very ordered file

aλφδ

βιτα

γδψηθ

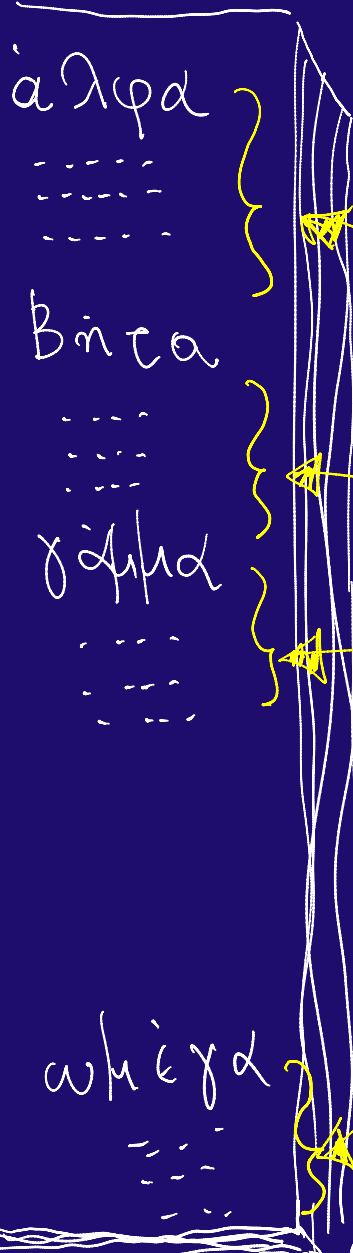
ωηιγδ

Intuitively, she expects:

words starting with a **given letter** = (all words)/24

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roughly equal
sizes

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ἀλφά

βίτα

γάμμα

ωμέγα

Intuitively, she expects:

words starting with a **given letter** = **(all words)/24**

So, she expects to find "αύγουστος" at:

1/24-th part of the file

Knuth's Example: a **lexicon** is a very nice and very ordered file

ἀλφα

βίτα

γάμμα

ωμέγα

roughly equal
sizes

!!!WARNING!!!

In Greek, there are too many ... "παπά*" words

Knuth's Example: a **lexicon** is a very nice and very ordered file

αλφα
.....
.....

βιτα
.....
.....

γίψα
.....
.....

παπά...
.....
.....

ωμεία
.....
.....

roughly equal
sizes

explodes

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In Greek, there are too many ... "παπά*" words

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α λ φ α

.....
.....

β η τ α

.....
.....

ζ α φ η α

.....
.....

παπά...

.....
.....

ω η ε γ α

.....
.....

roughly equal
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!!!WARNING!!!

In Greek, there are too many ... "παπά*" words

Hence, she **errs** by expecting "**ρομαντικός**" in the 17-th portion of the file, which deviates far to the end of file.

So, this is (a variant of) interpolation

Morals: **uniformity** matters (at least at present...)

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Knuth “**επικήρυξε**” the analysis of IS in his famous list of **10 most important problems in searching**. Thus, IS became the “**protagonist**” of at least **10 years** of intensive research.

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A great amount of papers focused on analyzing IS rigorously on **UNIFORM input distribution**:

Yao & Yao (1976),

G. Gonnet (1977)

Perl, Reingold (1977),

Perl, Itai, Avni (1978)

Gonnet, Rogers & George (1985)

They concluded to $\Theta(\log \log n)$ performance, **exponentially** better than Binary Search

What about **NON uniform** input distributions?

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D. Willard (1985),
Demaine, Jones & Patrascu (2004)

What about **Dynamic** files? (**insert/delete** keys)

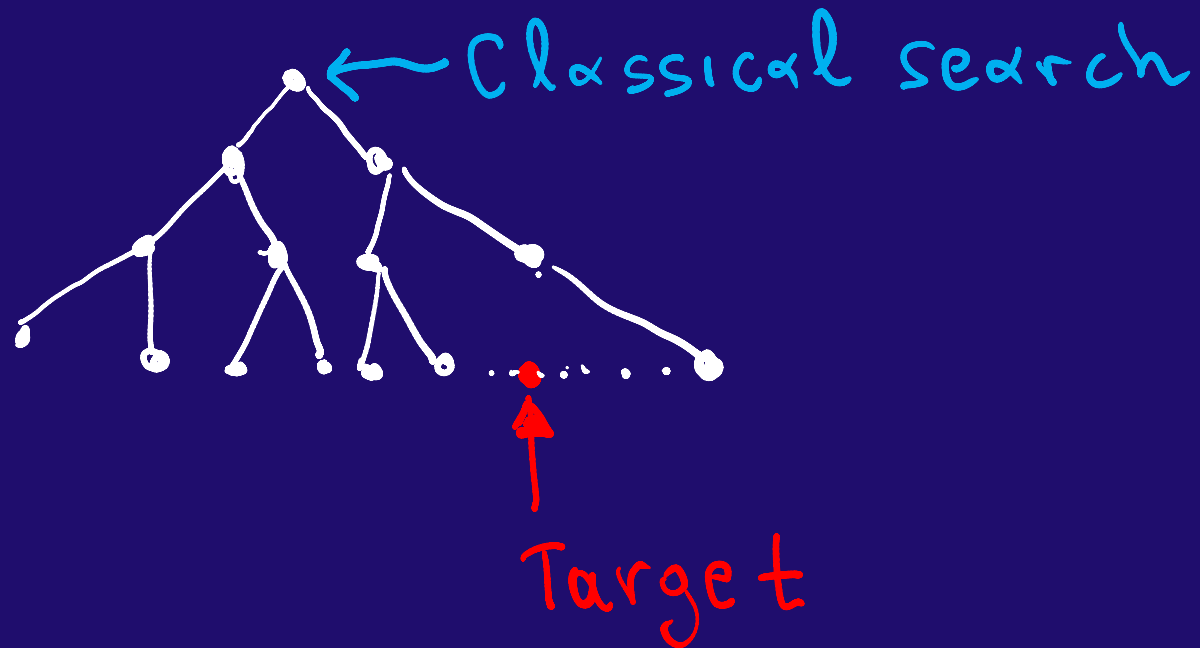
What about **Dynamic** files? (insert/delete keys)

Mehlhorn & Tsakalides (19??), also extended input distributions
Andersson & Mattson (1993), more extended input distributions

What about removing n from $\log\log(n)$?

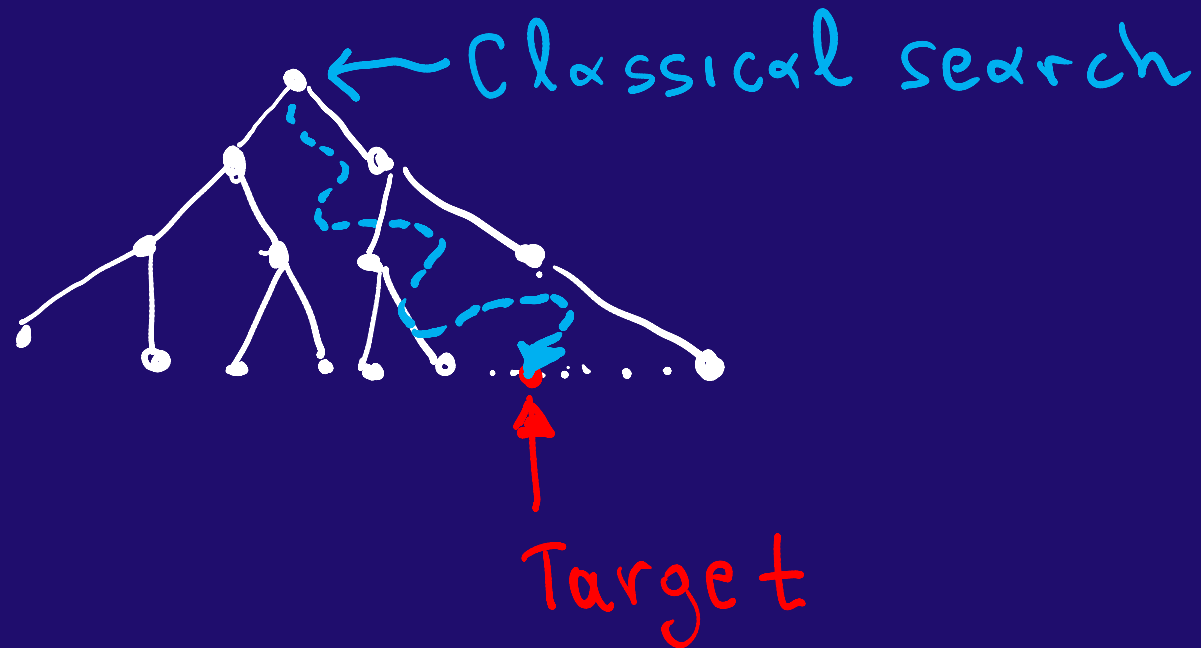
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Kaporis, Makris, Sioutas, Tsakalidis, Tsihlias & Zaroliagis (2003)
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 $d = \text{distance}(\text{target key, finger key})$



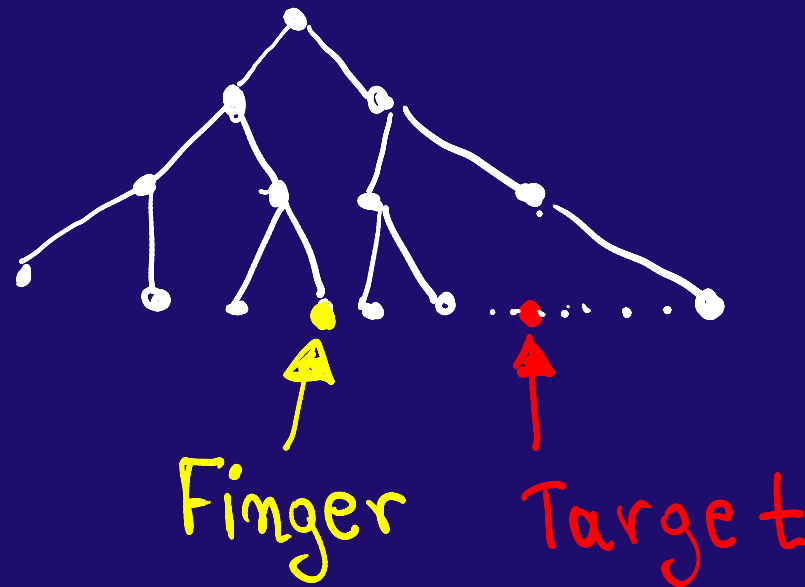
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Why not removing this "innocent" assumption???

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Why not removing this "**innocent**" assumption???

Because, strong experimental evidence showed that **all existing versions of IS** behave poorly on **finite** keys, e.g., **alphabetic tables**, Perl & Gabriel (1992)

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Why not removing this "**innocent**" assumption???

Because, strong experimental evidence showed that **all existing versions of IS** behave poorly on **finite** keys, e.g., **alphabetic tables**, Perl & Gabriel (1992)

Frankly, IS was **misled** by repetitions of **finite** keys
(recall: key repetition has **probability 0** in a continuous distribution)




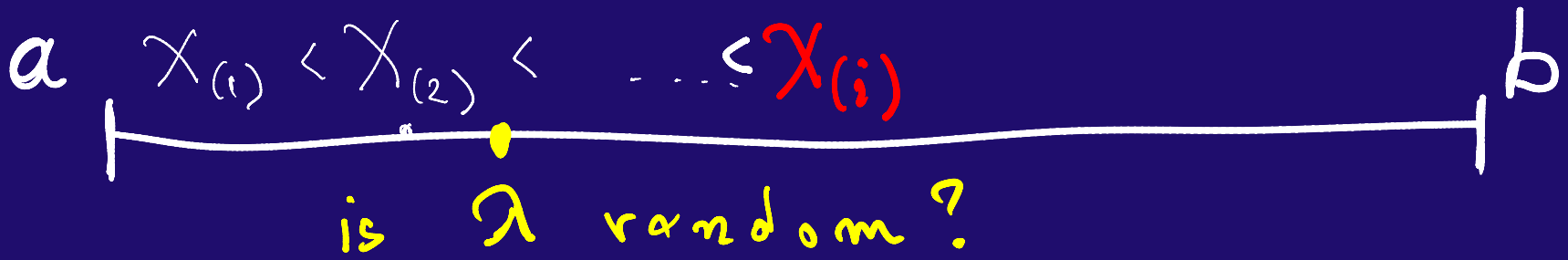
$$a \quad X_{(1)} < X_{(2)} < \dots < X_{(i)} < \dots < X_{(n)} \quad b$$

a $x_{(1)} < x_{(2)} < \dots < x_{(i)} < \dots < x_{(n)}$ b

$x_{(i)} > \text{target}$

a $X_{(1)} < X_{(2)} < \dots < X_{(i)}$ b





	λ	$P[X_{(2)} = \lambda \mid X_{(1)} = 3 \cap X_{(3)} = 10]$	$P[X = \lambda \mid 3 \leq X \leq 10]$
Analytic	3	0.00785	0.01481
Experimental	3	0.00755	0.01480
Analytic	4	0.04710	0.04445
Experimental	4	0.04707	0.04448
Analytic	5	0.10363	0.09779
Experimental	5	0.10364	0.09780
Analytic	6	0.17272	0.16299
Experimental	6	0.17311	0.16301
Analytic	7	0.22207	0.20956
Experimental	7	0.22099	0.20954
Analytic	8	0.22207	0.20956
Experimental	8	0.22228	0.20956
Analytic	9	0.17272	0.16299
Experimental	9	0.17280	0.16301
Analytic	10	0.05181	0.09779
Experimental	10	0.05252	0.09775