



# Fundamental Properties of Cryptocurrency in Distributed Systems

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Boston College



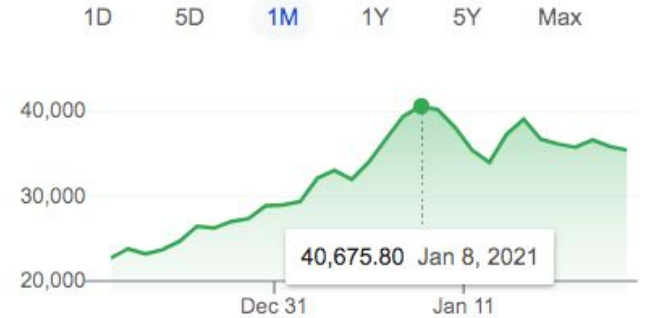
1 Bitcoin equals

**32,650.00 United States Dollar**

Jan 21, 7:45 PM UTC · Disclaimer

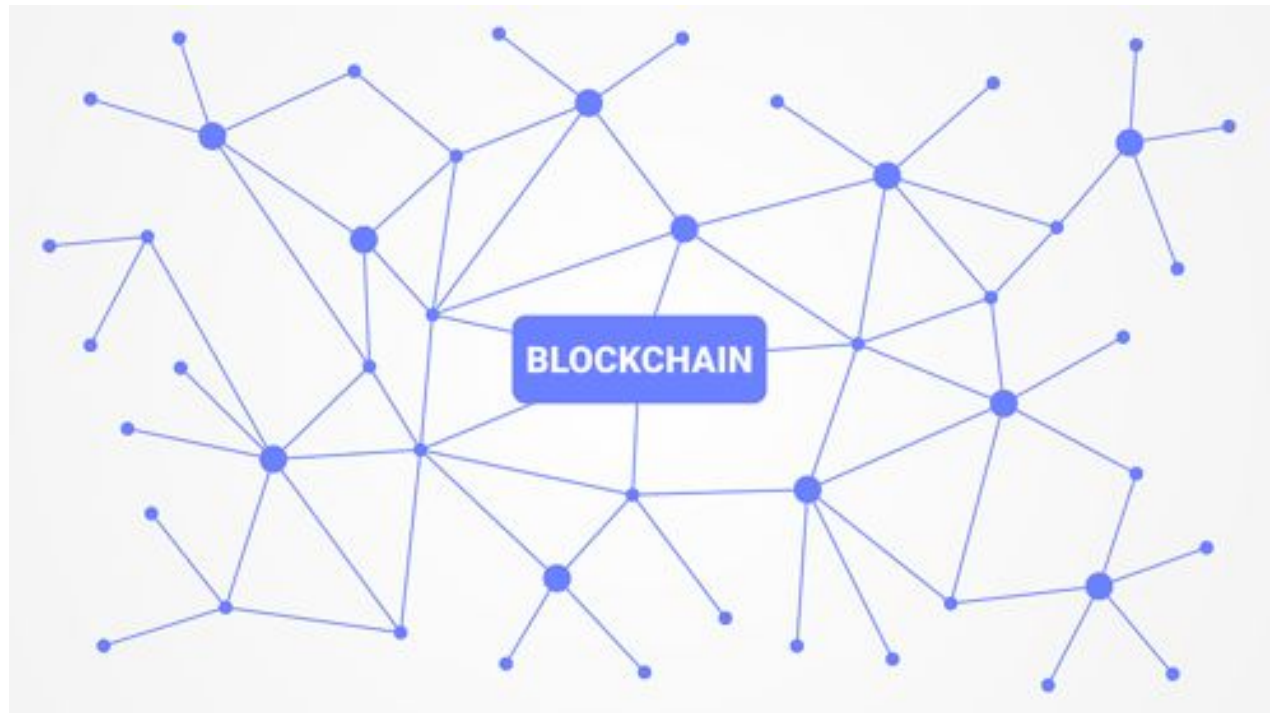
1	Bitcoin
32650.00	United States Dollar

Google search, 01/21/2021



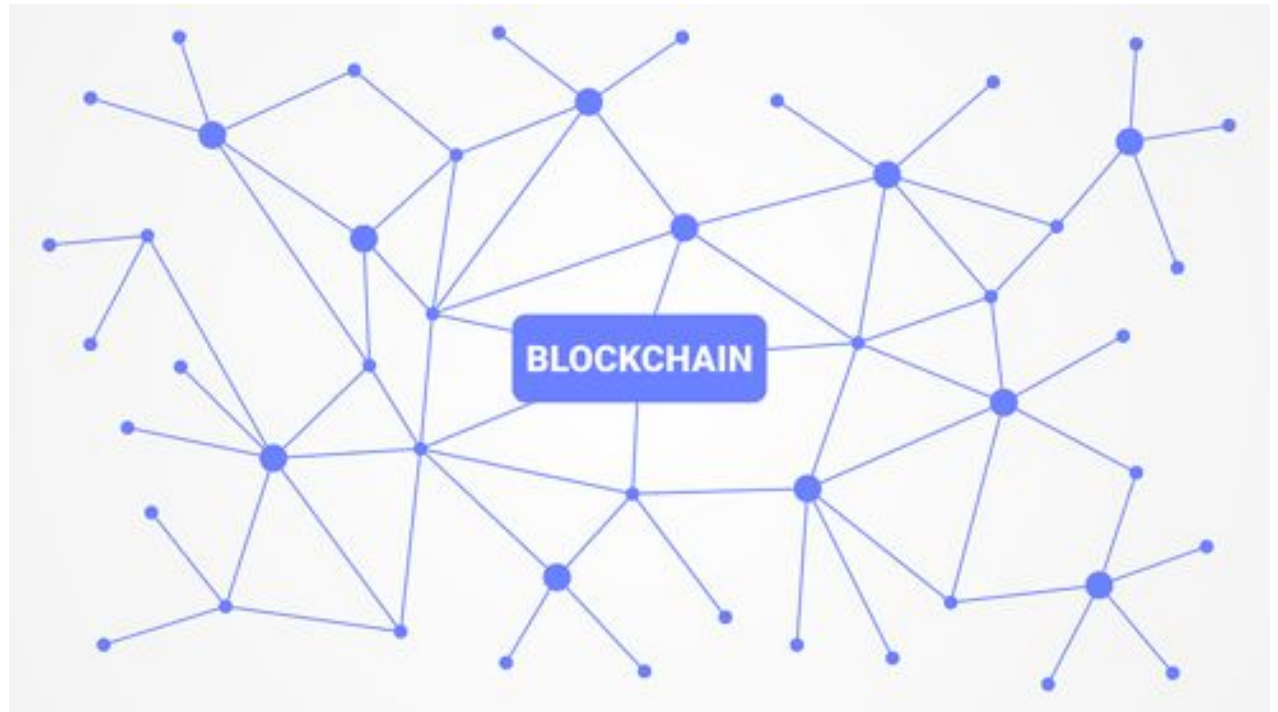
Data provided by Morningstar for Currency and Coinbase for Cryptocurrency

Decentralized  
Fault-tolerant  
Secure



ways of sharing and storing information

Decentralized  
Fault-tolerant  
Secure  
... many other



ways of sharing and storing information

# Fundamental Questions

## Blockchains and the Future of Distributed Computing

[Herlihy PODC 17]

- No formal abstraction of these objects has been proposed

## Formalizing and Implementing Distributed Ledger Objects

[Anta et al. NETYS 18]

- What is the service that must be provided by a distributed ledger?
- What properties a distributed ledger must satisfy?
- What are the assumptions made by the protocols and algorithms on the underlying system?
- Does a distributed ledger require a linked cryptocurrency?

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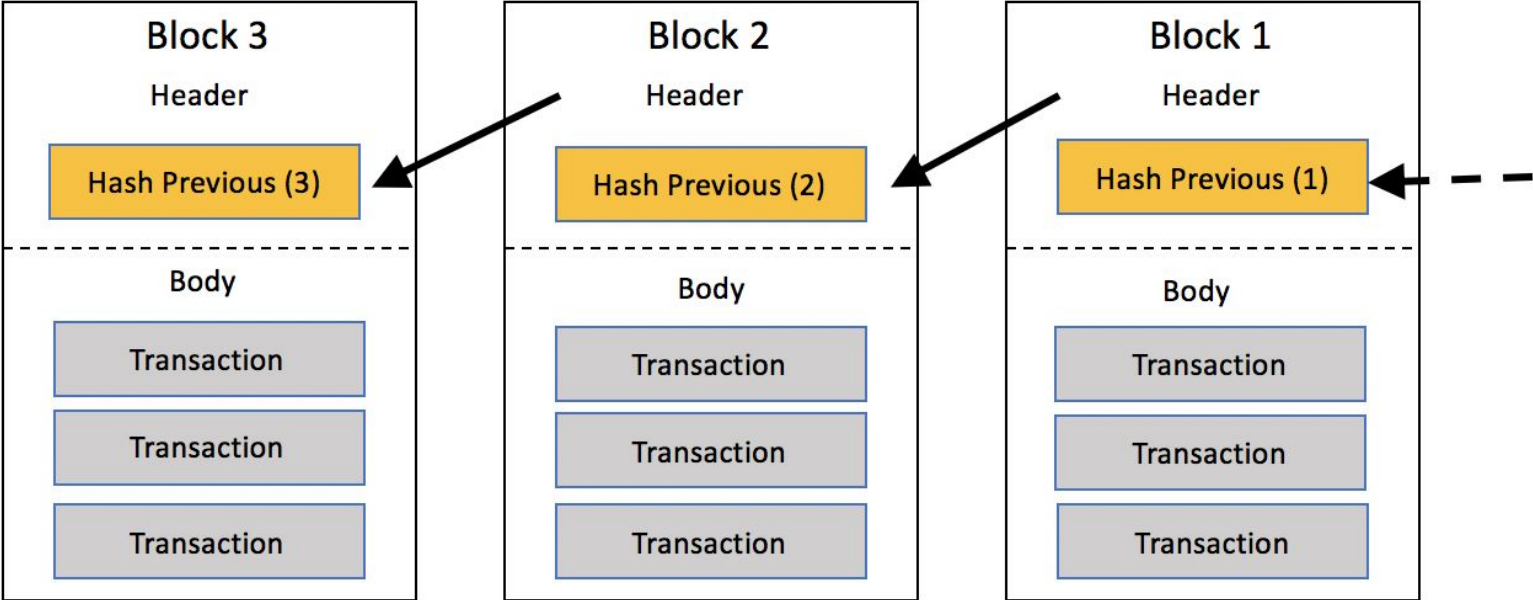
## Formalizing and Implementing Distributed Ledger Objects

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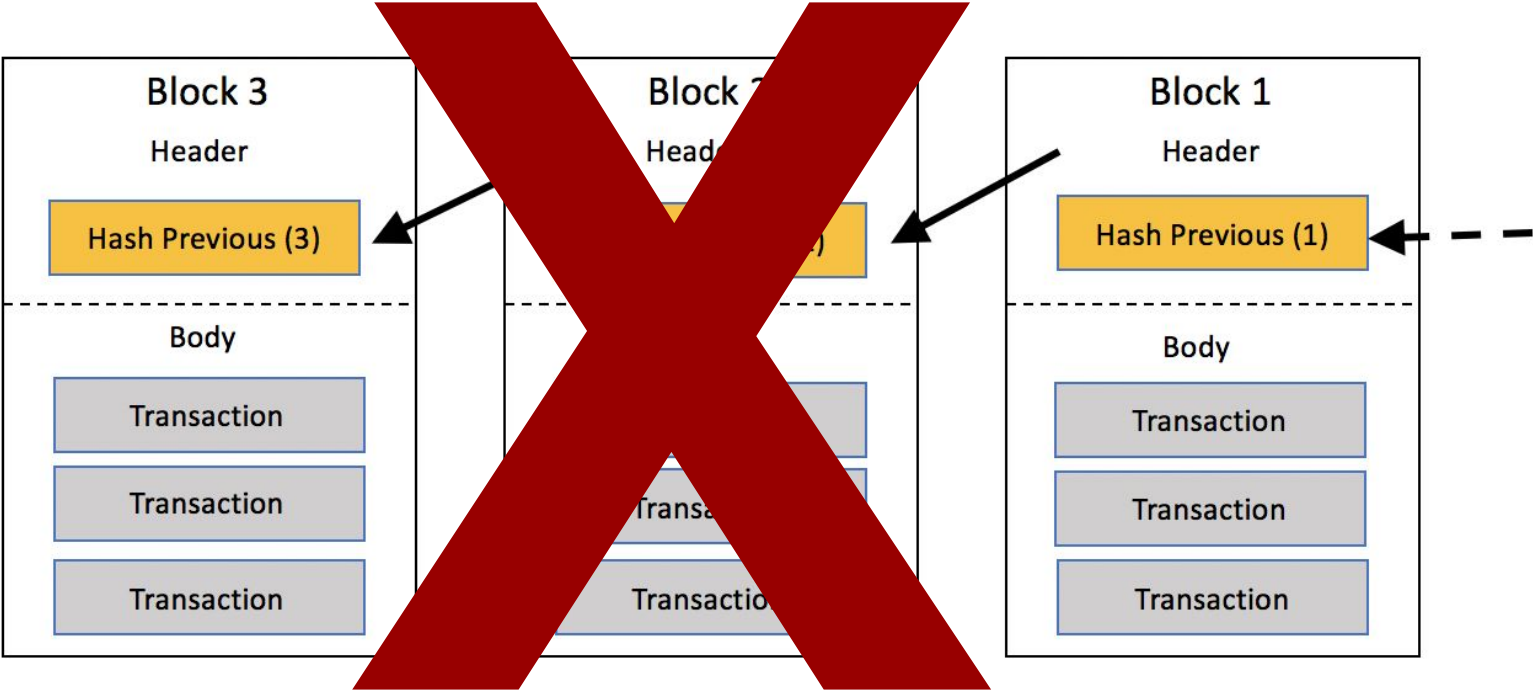
# Why Blockchain is Hard?

# Blockchain?



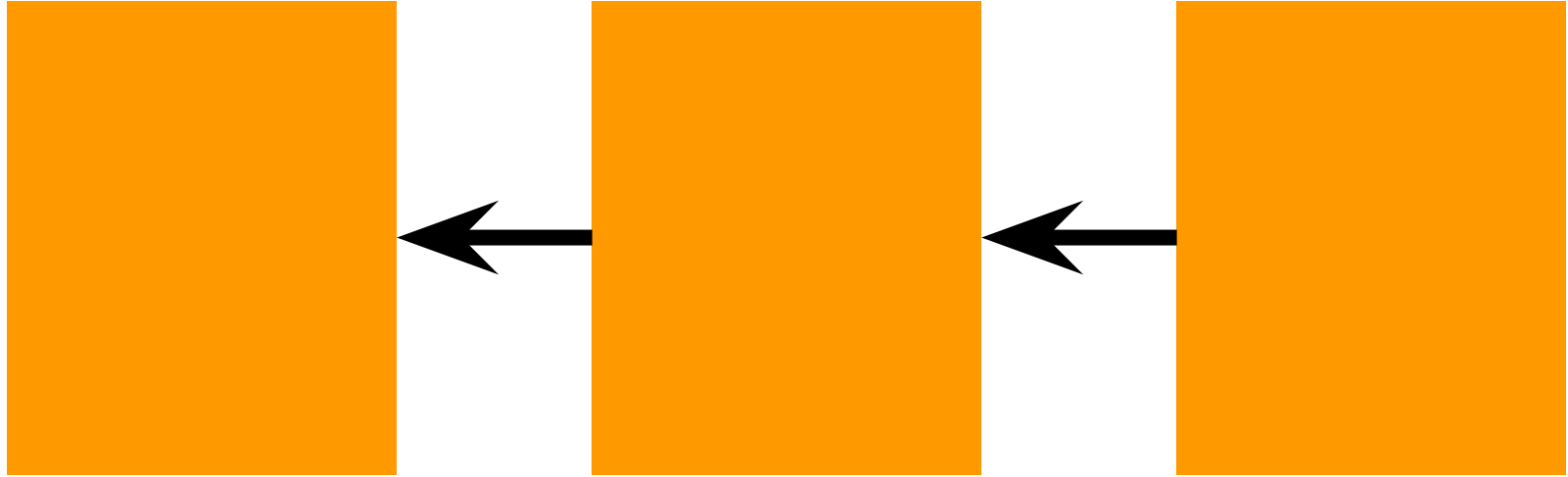


# Blockchain?



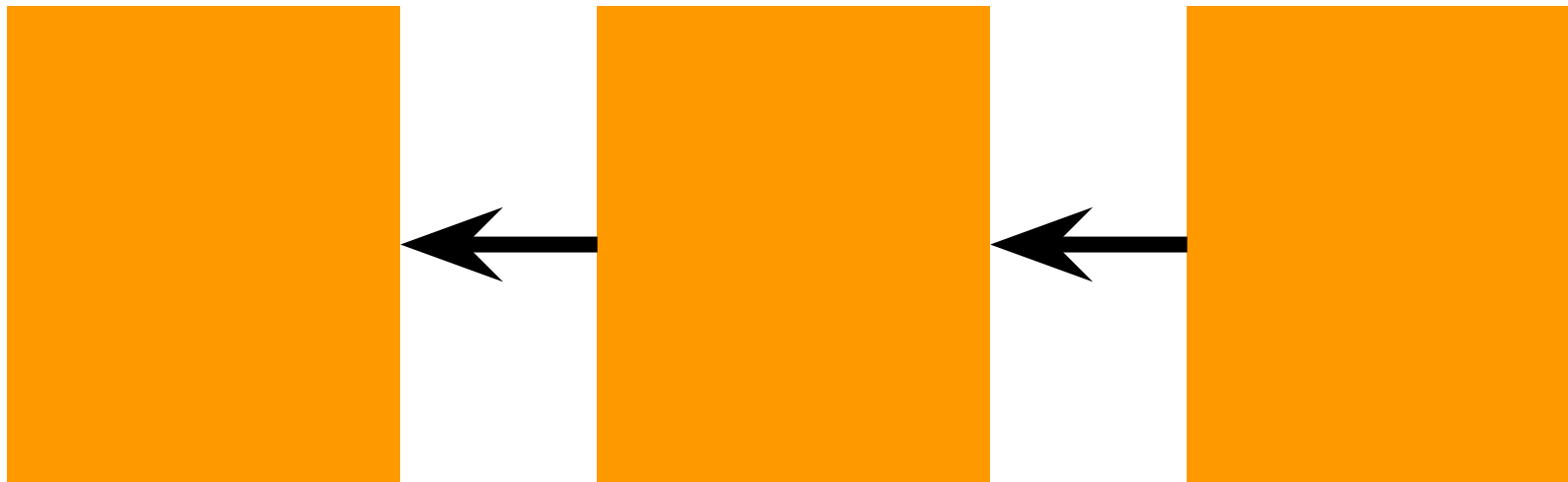
Blockchain?

Block + Chain!



# Blockchain?

# Block + Chain!

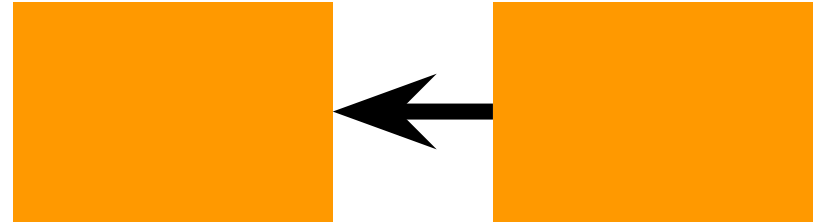


In this talk, no crypto detail

- What Block?
- How to link?
- How secure?
- How to mine?

# Block + Chain!

[in Distributed Computing]



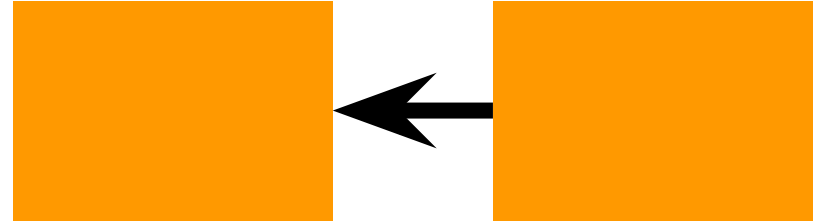
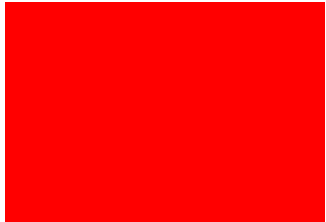
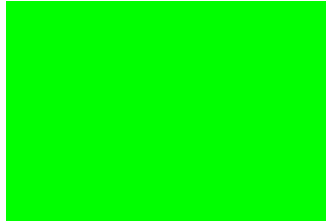
Bitcoin on a high-level:

In each round,

- Nodes exchange blocks (mining)
- Nodes “agree on” a block

# Block + Chain!

[in Distributed Computing]



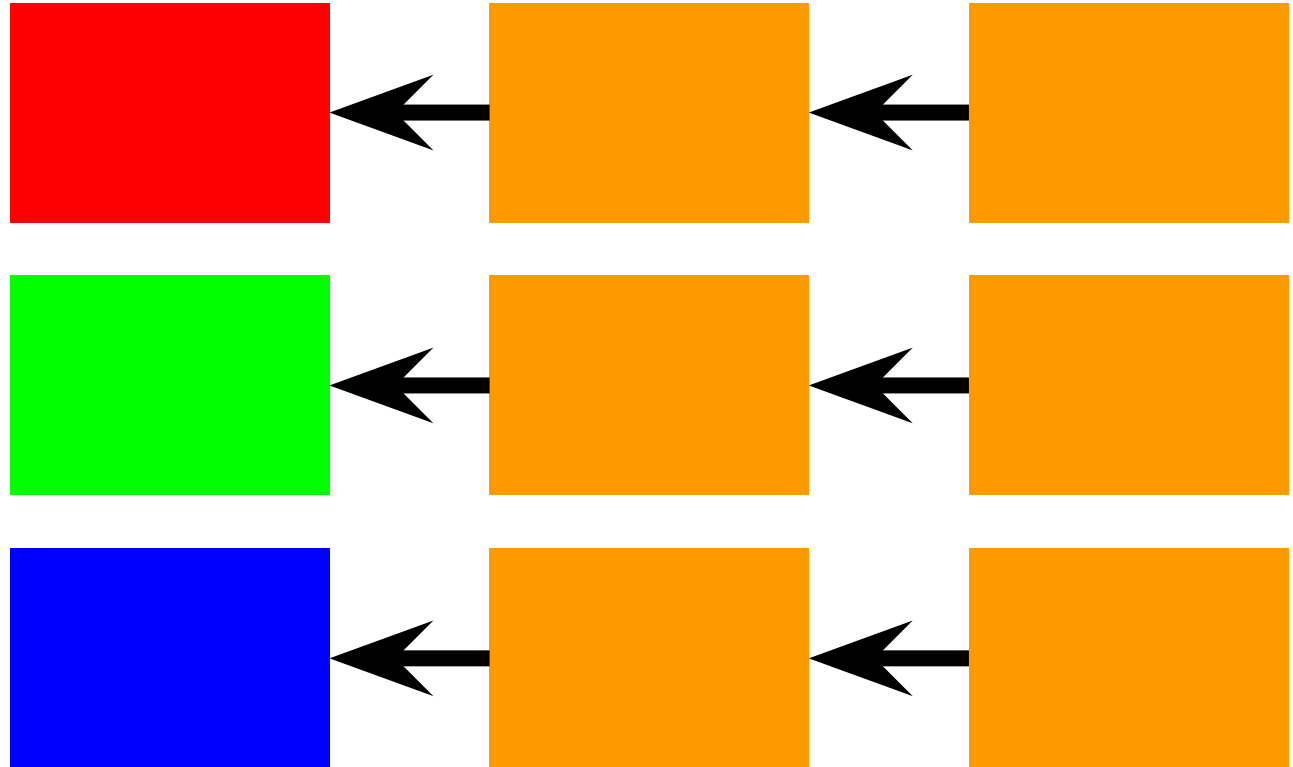
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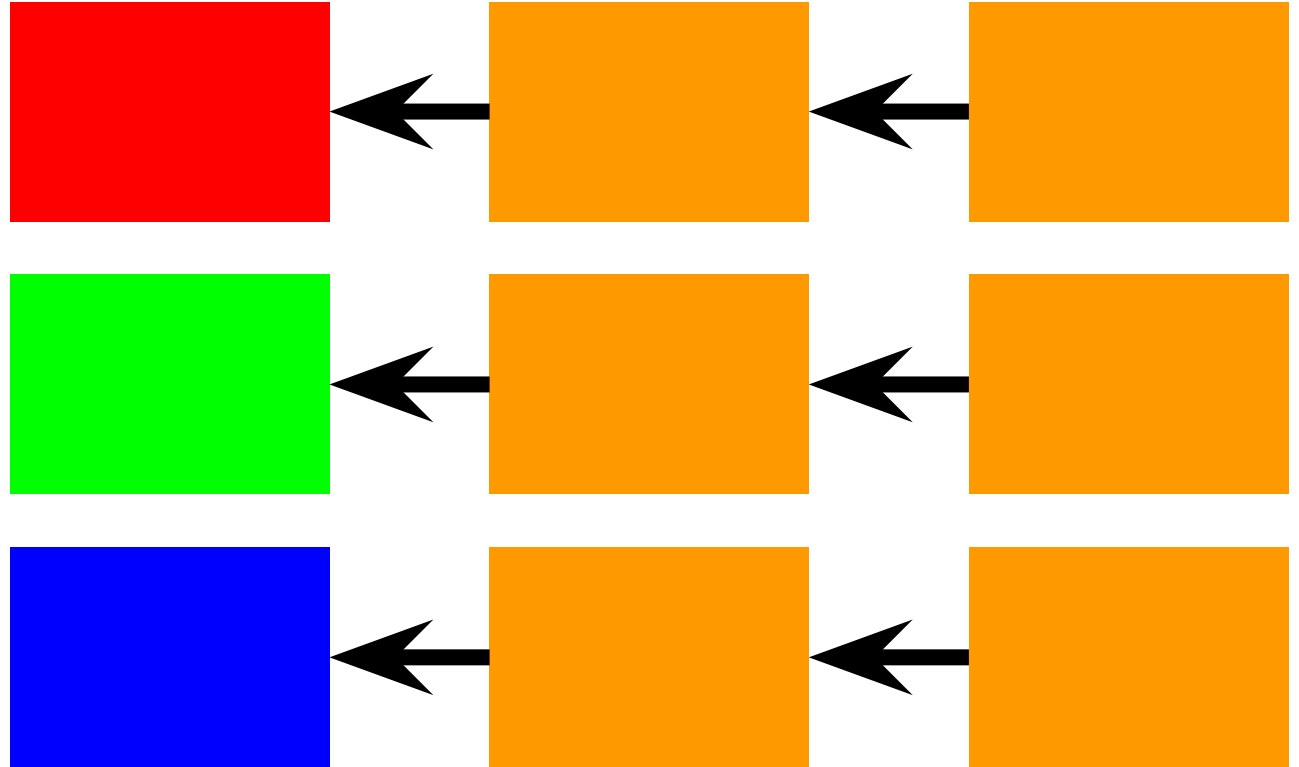
[in Distributed Computing]



# Block + Chain!

[in Distributed Computing]

**disagreement** ⇒  
double-spending  
attack



# Block + Chain!

[in Distributed Computing]

**disagreement**  $\Rightarrow$   
double-spending  
attack



**FLP Result**  
**[JACM 85]:**  
Fault + Async. +  
Consensus  
= Impossible





Most common approach:  
Proof-of-XXX

# Is Consensus necessary?

Is Consensus necessary?  
No

# The Consensus Number of a Cryptocurrency

Rachid Guerraoui, Petr Kuznetsov, Matteo Monti, Matej Pavlovic, Dragos-Adrian Seredinschi [PODC 2019]

<https://arxiv.org/pdf/1906.05574.pdf>

# Model

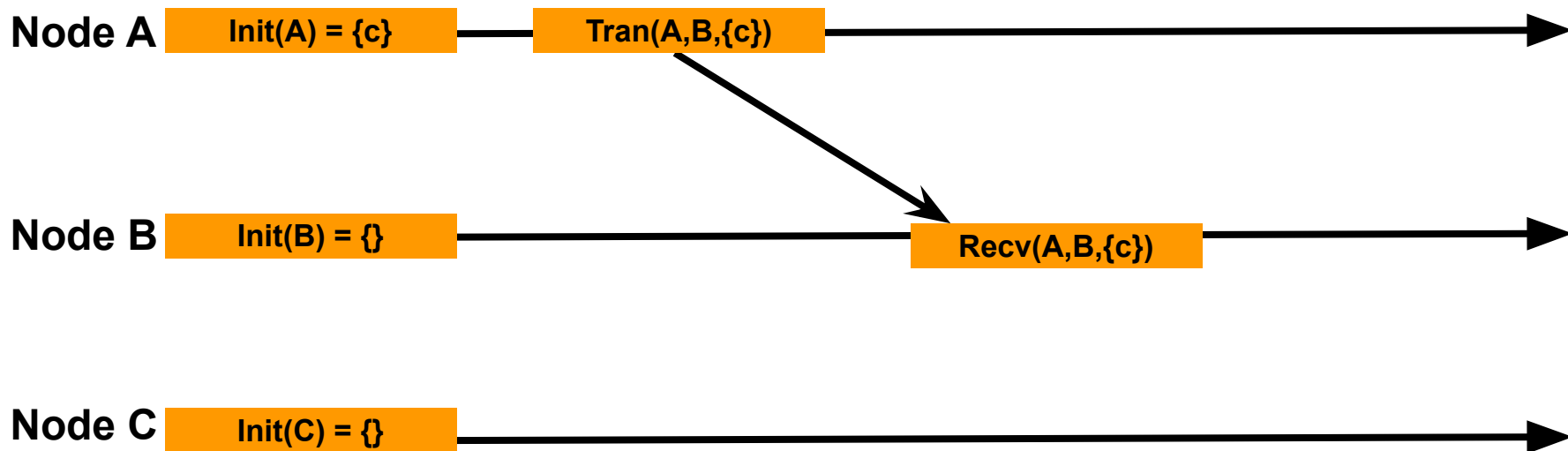
Asynchronous network:	arbitrary message delay
Permissioned, static system:	a fixed set of nodes $[1, \dots, n]$
Crash fault:	up to $f$ fail-stop failures

# Cryptocurrency: Working Definition

A cryptocurrency is a virtual asset that relies on cryptography tools to prevent counterfeit or double-spend.

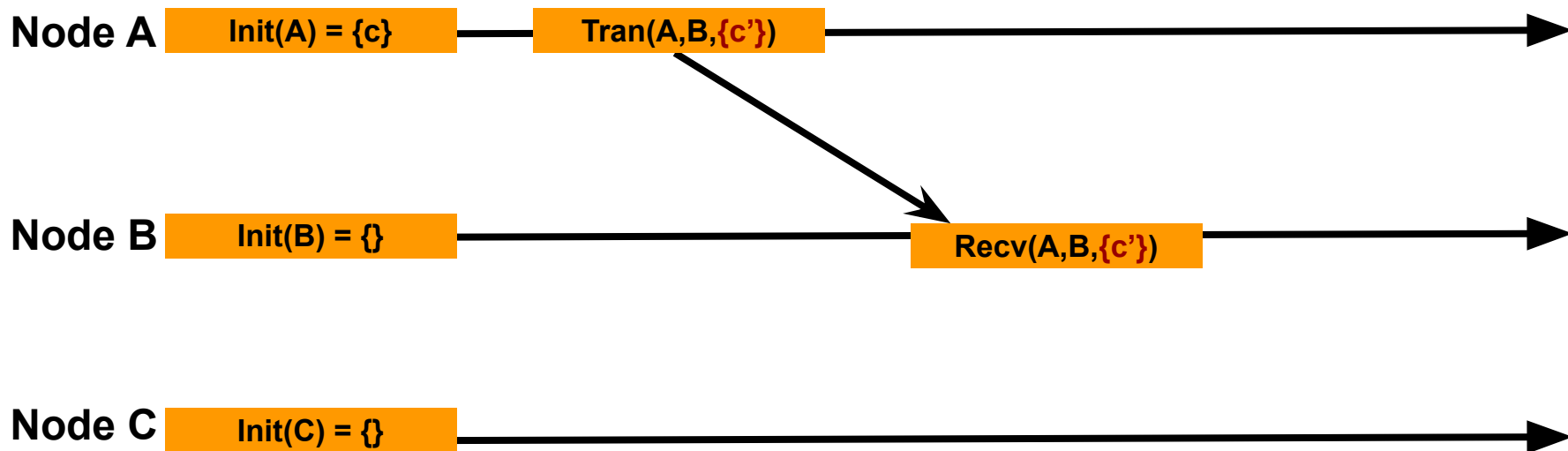
# Cryptocurrency: Abstraction

A cryptocurrency is a virtual asset that relies on cryptography tools to prevent counterfeit or double-spend.



# Cryptocurrency: Counterfeit

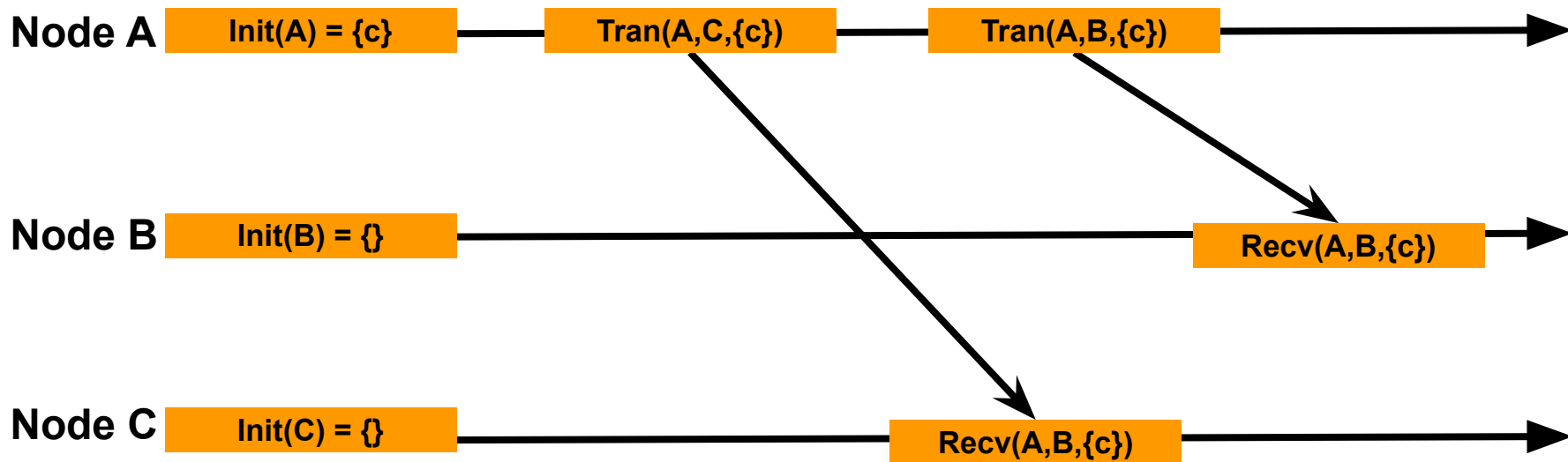
A cryptocurrency is a virtual asset that relies on cryptography tools to prevent **counterfeit** or double-spend.





# Cryptocurrency: Double-spend

A cryptocurrency is a virtual asset that relies on cryptography tools to prevent counterfeit or **double-spend**.



# What is the (concurrent) data structure?

## Asset transfer object [PODC 2019]

- Each node has an account
- ATO state: balance of each account
- Transfer(A,B,x):
  - Decrease A's account by x
  - Increase B's account by x
- Read(A): return A's balance

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**No overdraft**

# Atomic Snapshot Object [Afek et al. JACM 93]

Object partitioned into  $n$  segments

Each segment is “owned” by a node (single-writer)

**Update:** write a value to own segment

**Scan:** read values from all segments -- take a snapshot

Operations **linearizable** -- a total order that follows the real-time order

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ID	1	2	3
Value			

# ASO-based Cryptocurrency [PODC 2019]

## Intuition:

- $i$ 's entry = all outgoing transfers at node  $i$

ID	1	2	3
Value	{(1, 2, 100)}		

# ASO-based Cryptocurrency [PODC 2019]

Read(A):

- $S \leftarrow AS.scan$
- Return A's balance in S

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	Balance
1	0
2	50
3	50

ID	1	2	3
Value	{(1, 2, 100)}	{(2, 3, 50)}	

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Read(A):

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 $OP[A] \leftarrow OP[A] \cup \{(A,B,x)\}$   
 $AS.update(OP[A])$

ID	1	2	3
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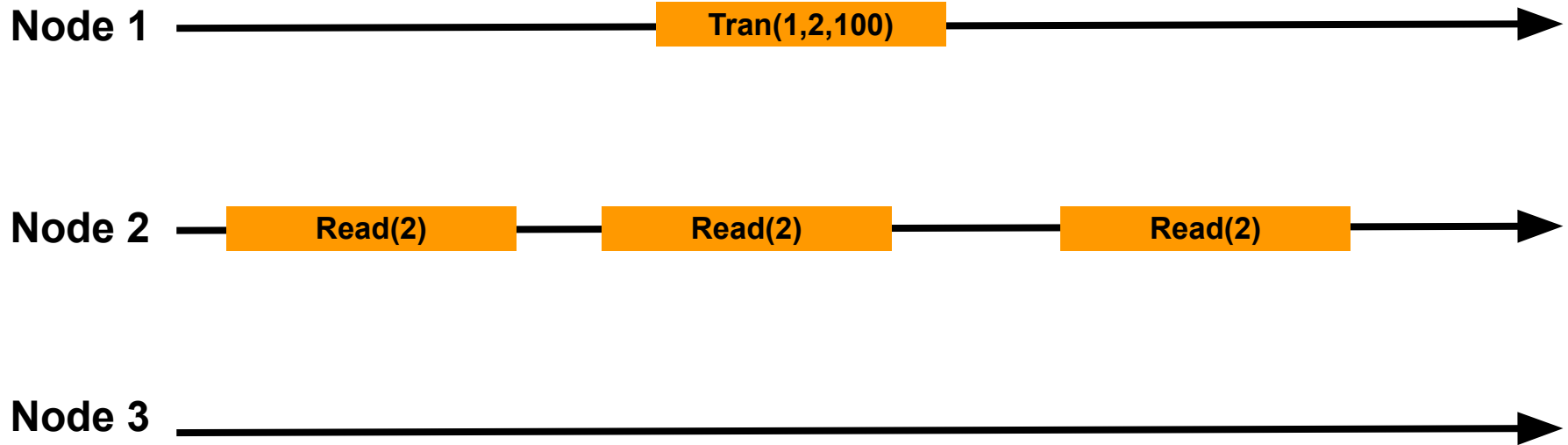
Transfer(A,B,x):

learn new incoming tx's

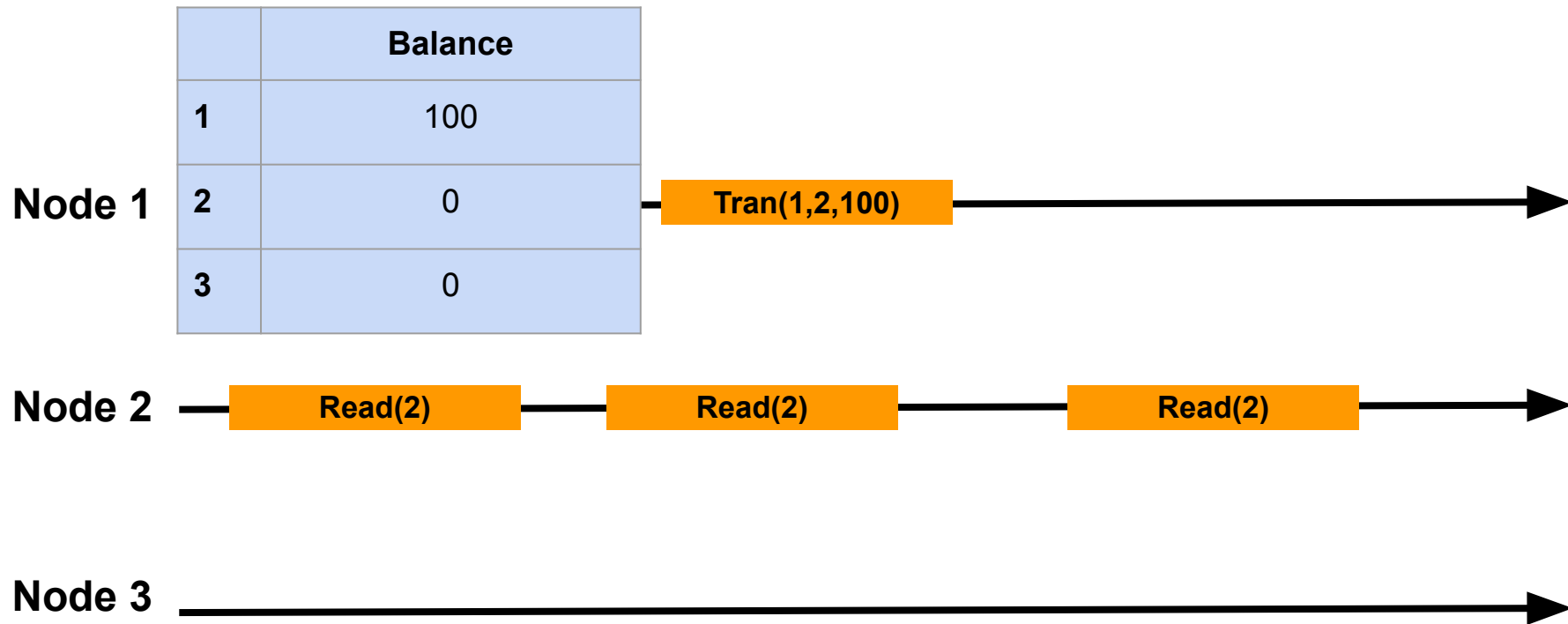
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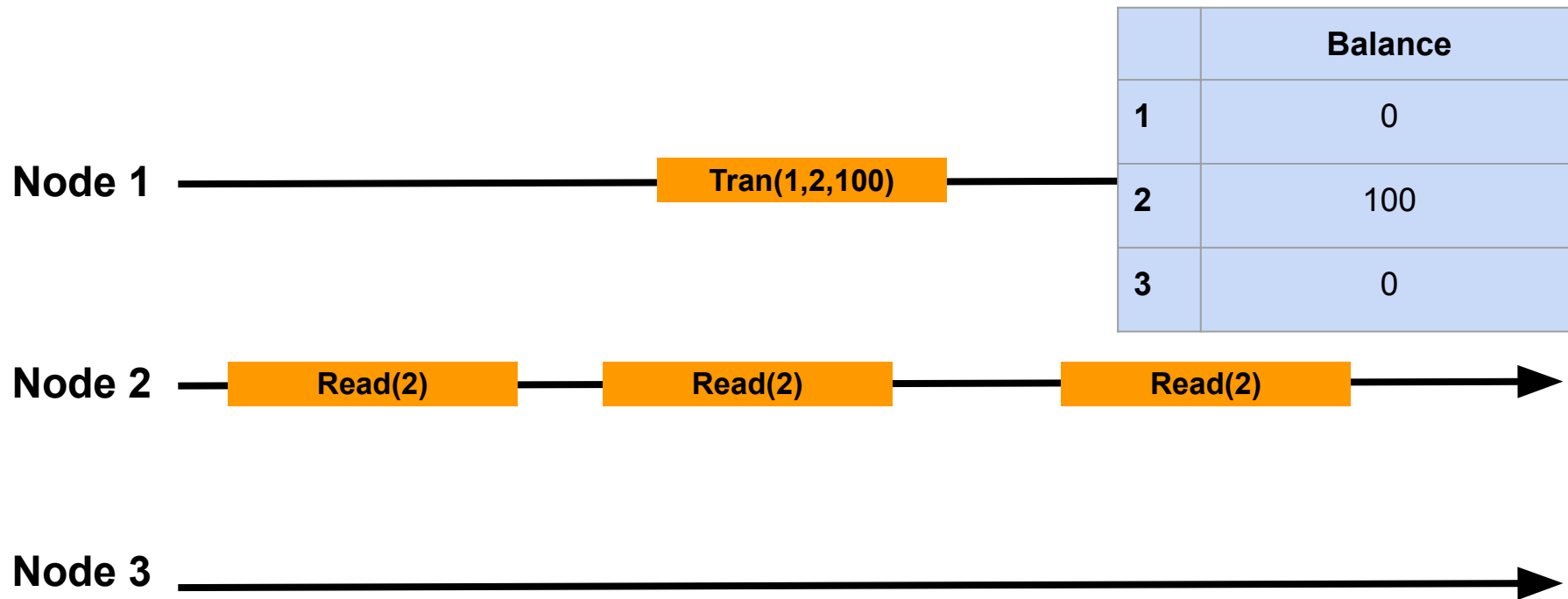
# Example



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1	0
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no overdraft

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- If valid transaction:  
 $OP[A] \leftarrow OP[A] \cup \{(A,B,x)\}$   
 $AS.update(OP[A])$

Update A's outgoing tx's

ID	1	2	3
Value	$\{(1, 2, 100)\}$	$\{(2, 3, 50)\}$	

# Example

	Balance
1	0
2	100
3	0

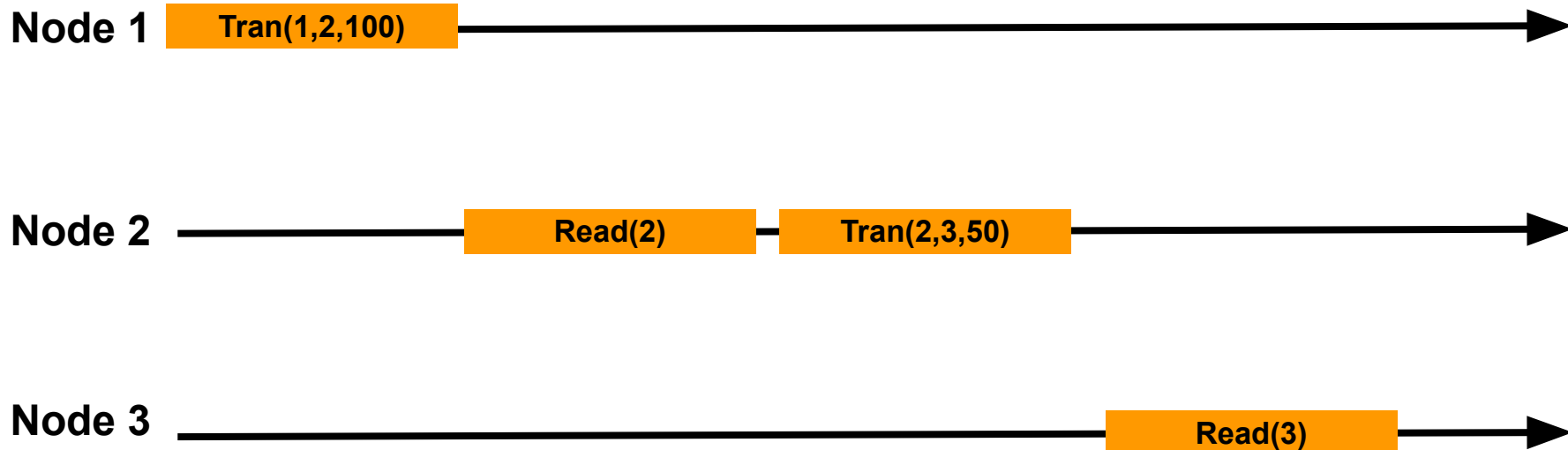
Node 1 **Tran(1,2,100)** →

Node 2 → **Read(2)** →

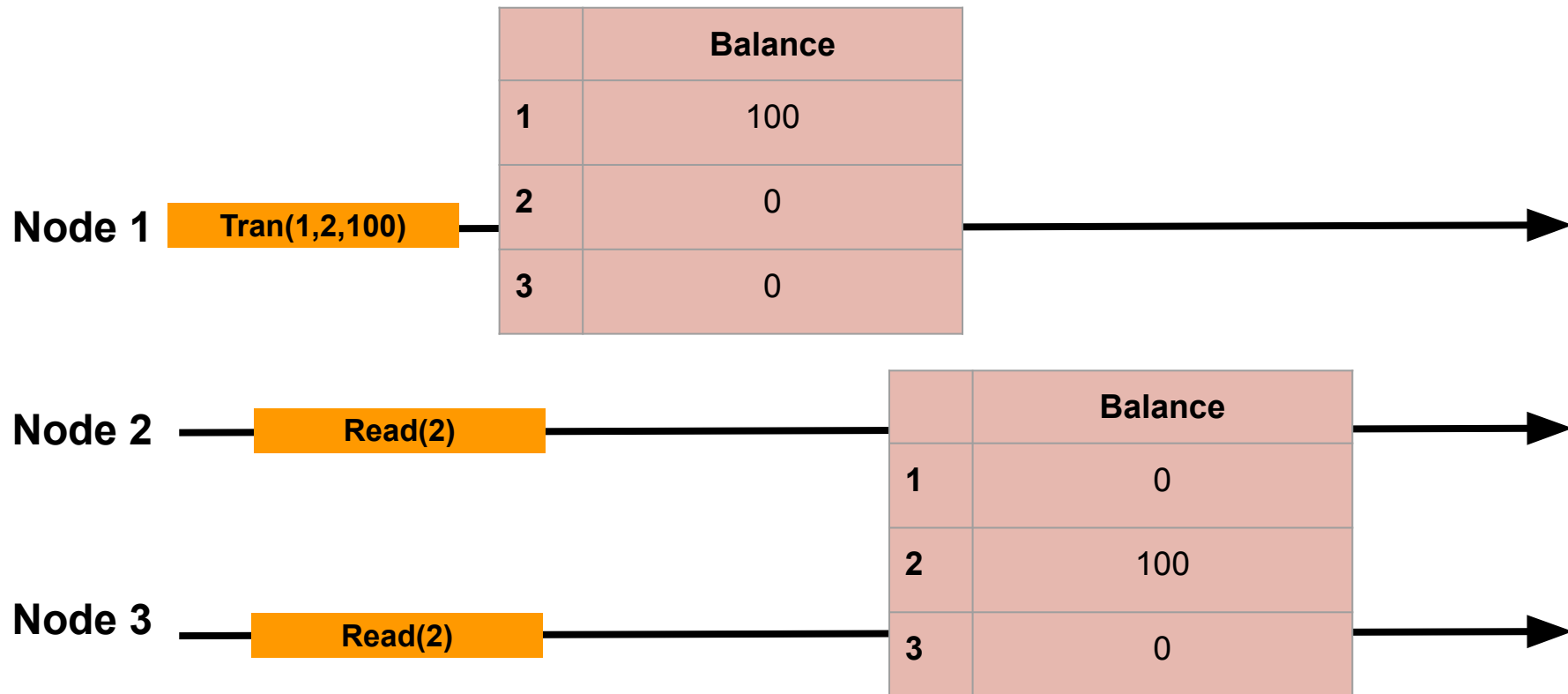
Node 3 →

# Example

	Balance
1	0
2	50
3	50



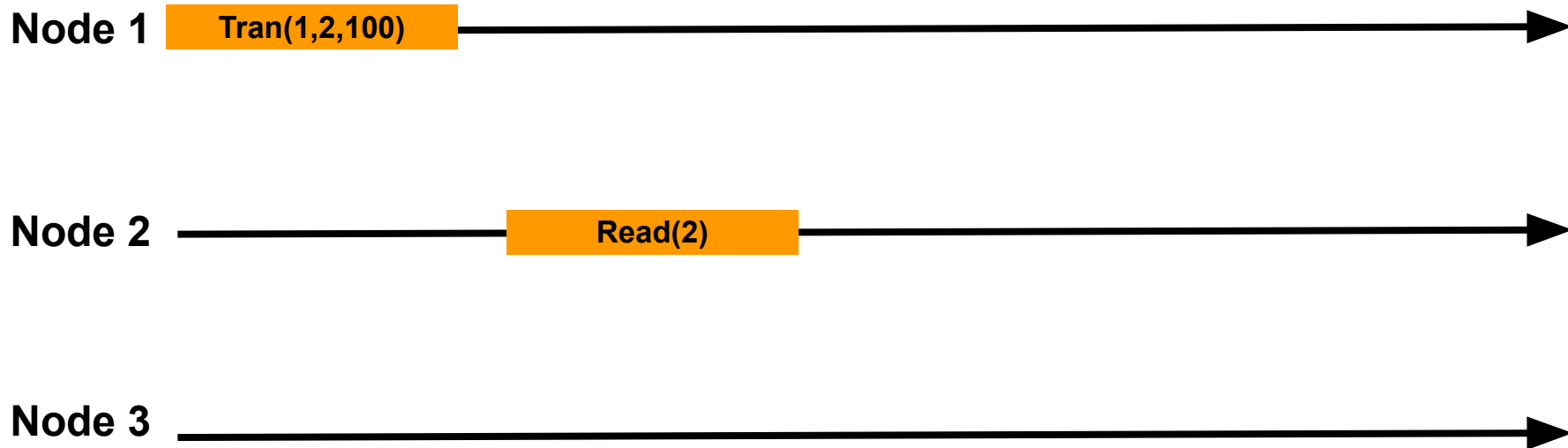
# Is this possible?



# Linearizability:

Total order + **Real-time order**

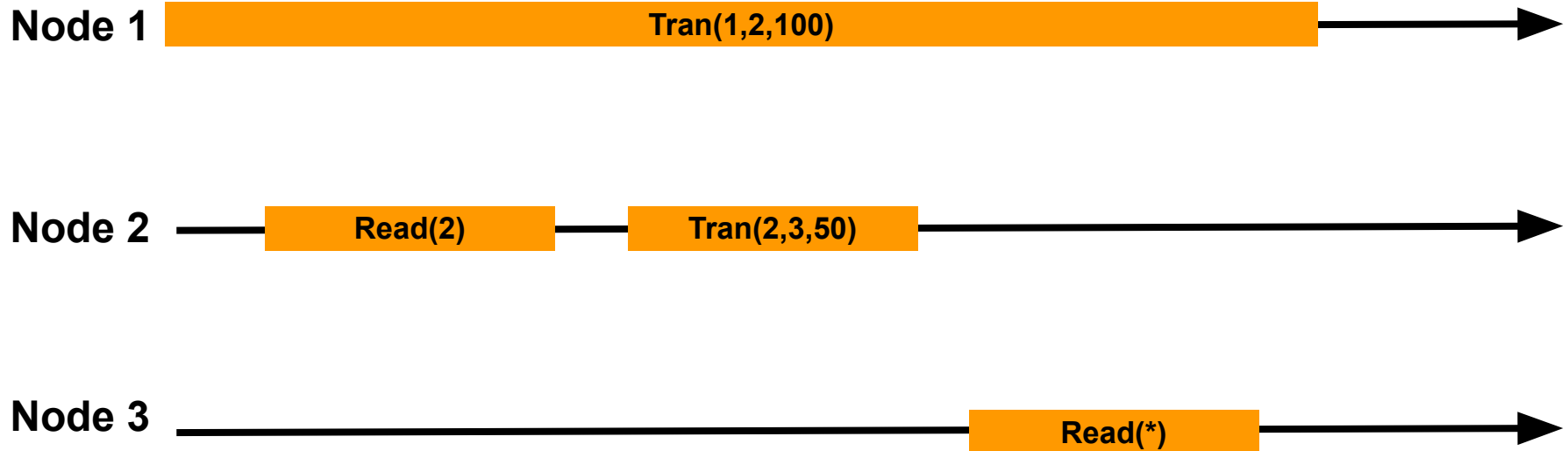
	Balance
1	0
2	100
3	0



# Linearizability:

**Total order** + Real-time order

	Balance
1	0
2	50
3	50



So what?

ASO can be implemented in asynchronous systems!



It's all good, but ...

Is it scalable and highly available?

CAP Theorem [Brewer PODC 00, Gilbert/Lynch 02]

**Consistency:** right response to each request

**Availability:** termination eventually

**Partition tolerance:**

unreliable comm. network

CAP Theorem [Brewer PODC 00, Gilbert/Lynch 02]

**Consistency:** right Impossible to have all three!

**Availability:** term

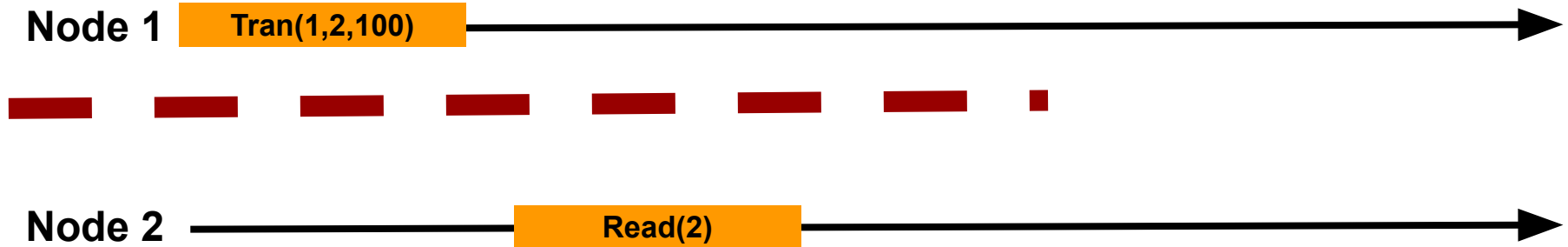
**Partition tolerance:**

unre

When there is a partition,  
choose consistency or  
availability

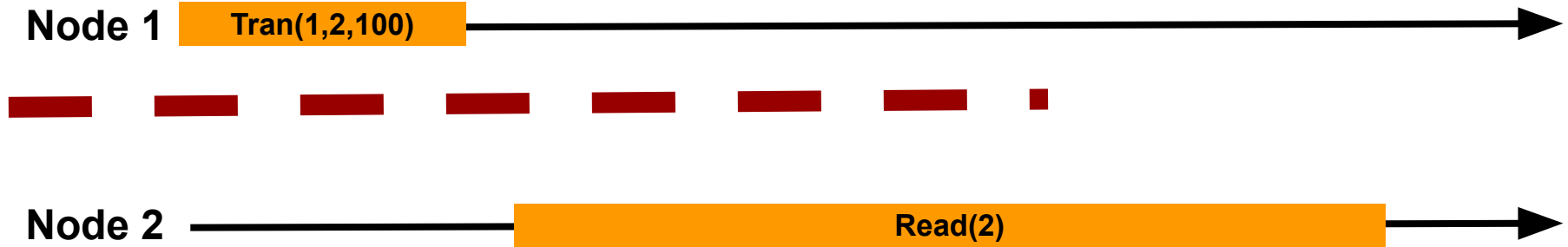
# ASO-Crypto: Incorrect

	Balance
1	100
2	0
3	0



# ASO-Crypto: Slow

	Balance
1	0
2	100
3	0



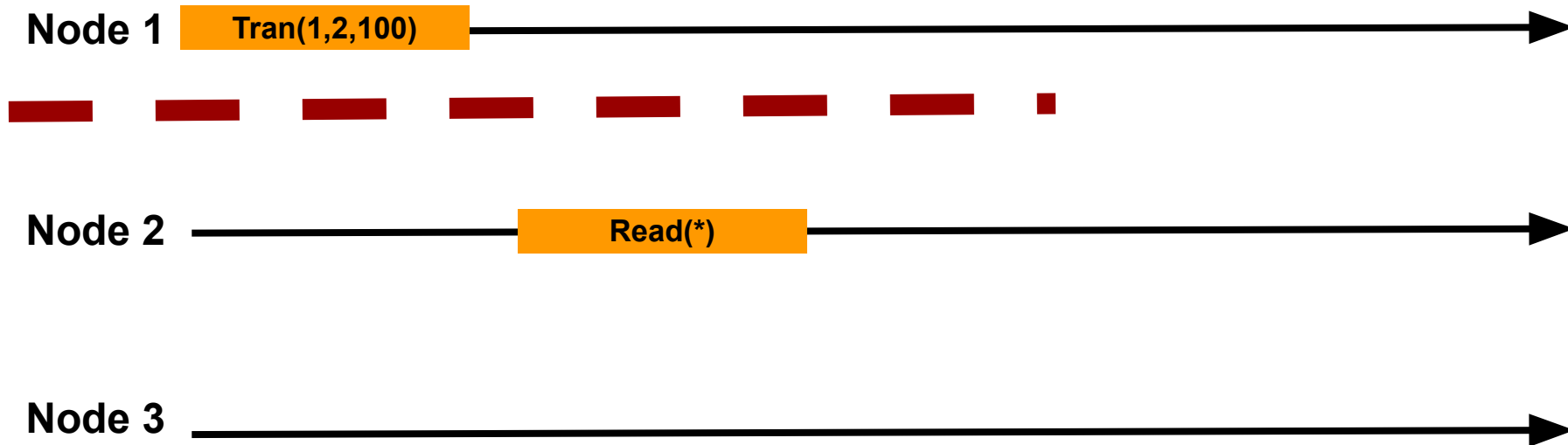
# Cryptocurrency in a Partitioned Network? (under submission)

**Key observation:**  
Pending transactions  
(Delivered but not applied)

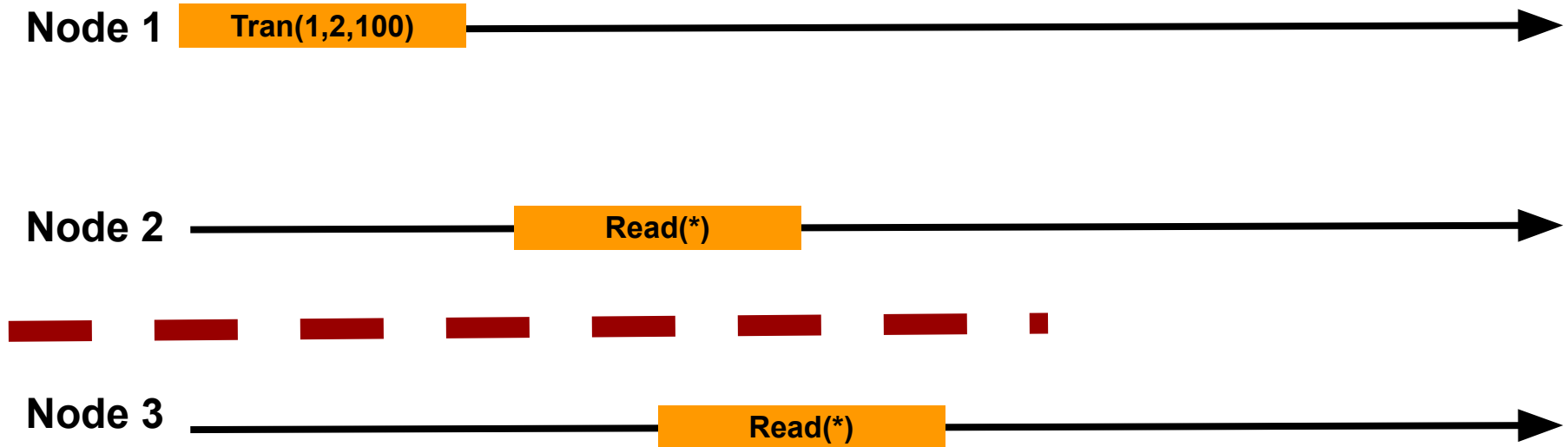


# Modern Bank

	Balance
1	100
2	0
3	0



# Modern Bank



# Modern Bank

	Balance
1	0
2	100
3	0

Node 1

Tran(1,2,100)



Node 2

Read(\*)

	Balance
1	100
2	0
3	0



Node 3

Read(\*)

**Key tool:**  
abstracting consistency guarantees

# Operations

## Transfer

## Read

**Audit:** return the “validity proof” of all the outgoing transactions

**Valid transaction:** no double-spend, no counterfeit, no overdraft

# Properties

- Eventual delivery:** Tx from the same partition is eventually applied
- Local operation:** No communication needed to complete an operation
- Read-my-write:** Read reflects the effect of all the prior outgoing transfers
- Auditability:** One is able to present validity proof
- Validity:** All applied transactions are valid

# Properties under CAP framework

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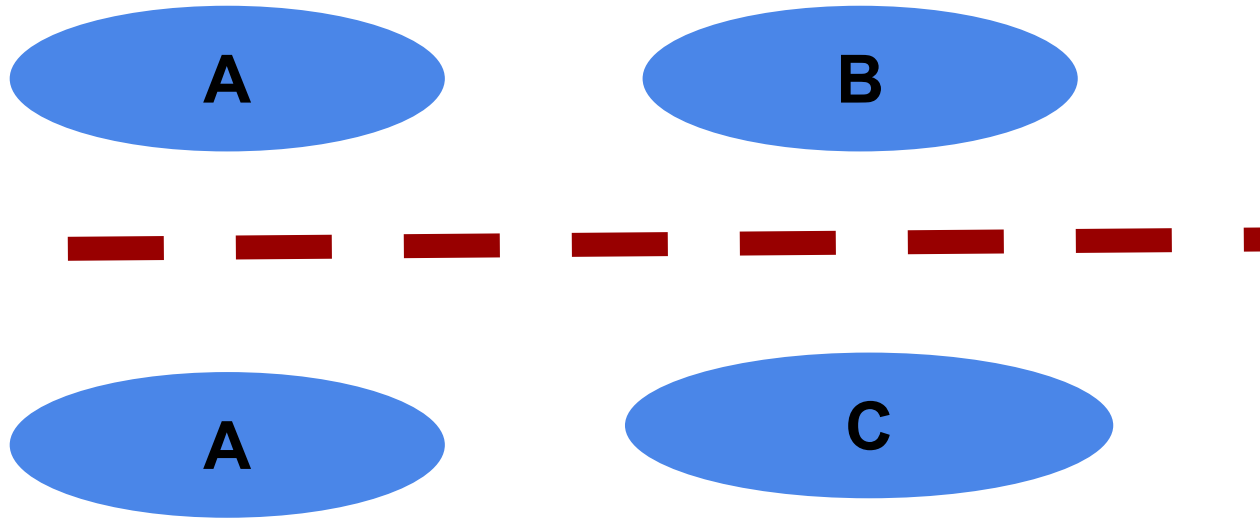
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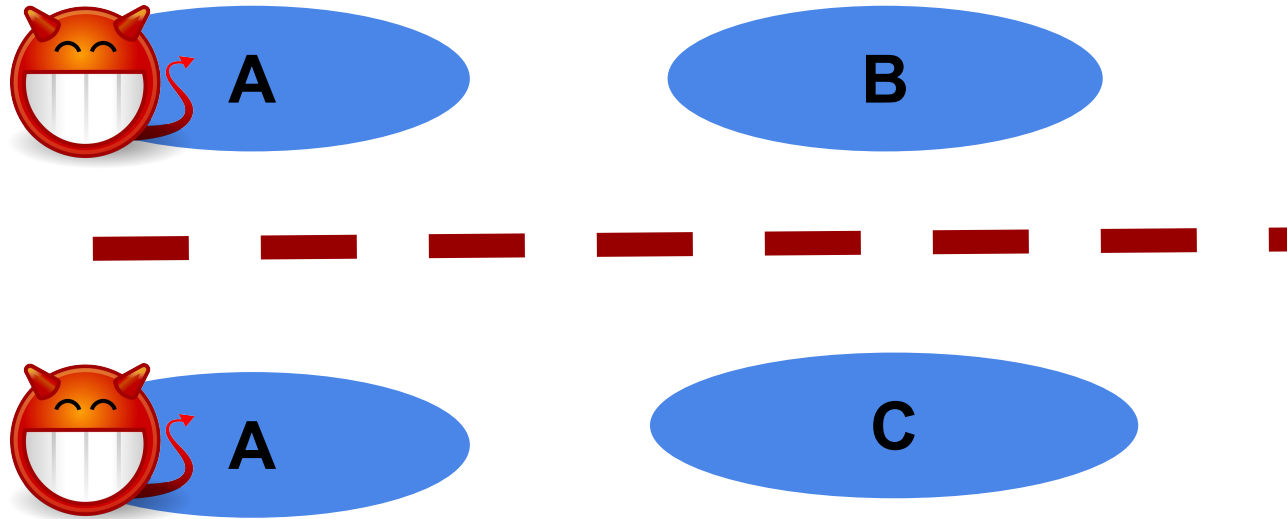
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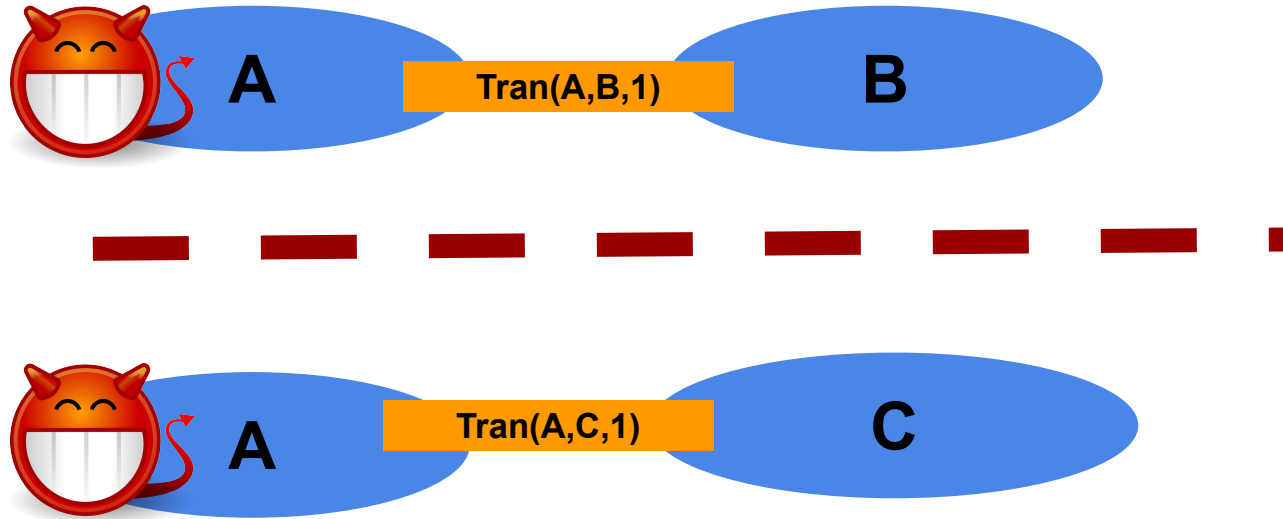
Thm1: Byzantine node + Eventual delivery +  
Partition-tolerance = impossible



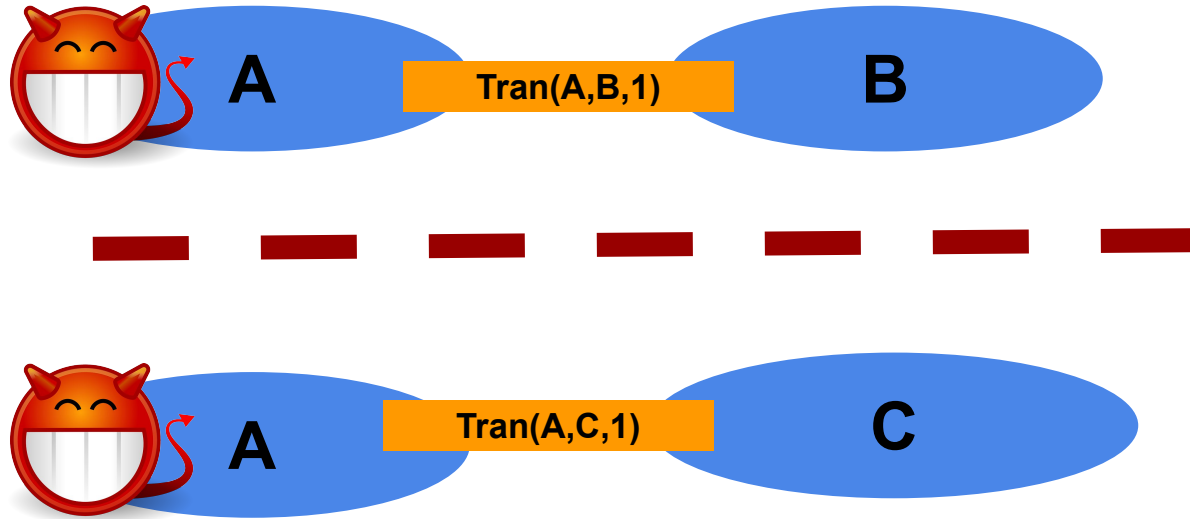
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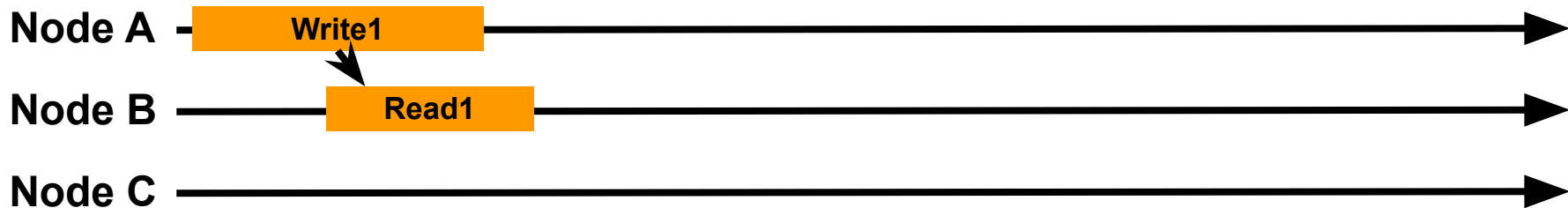
**Don't apply:**  
violating eventual  
delivery

**Apply:**  
double-spend

# Causal consistency

[Ahamad et al. DC 95]

On a high-level, causal consistency lets each node observe the entire causal history (happens-before relation)

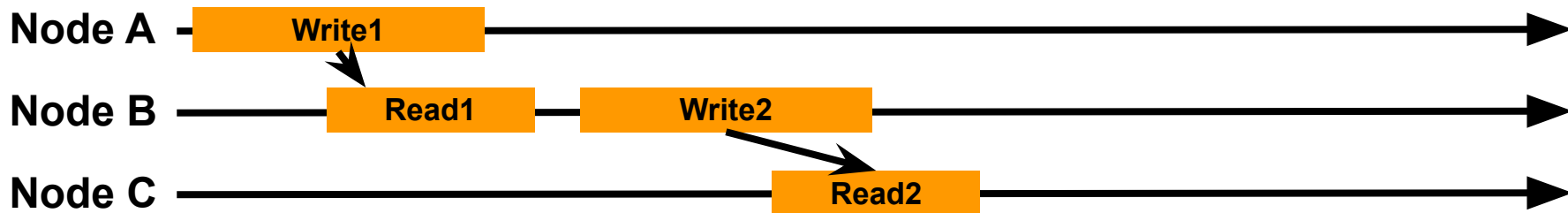


Write1 → Read1

# Causal consistency

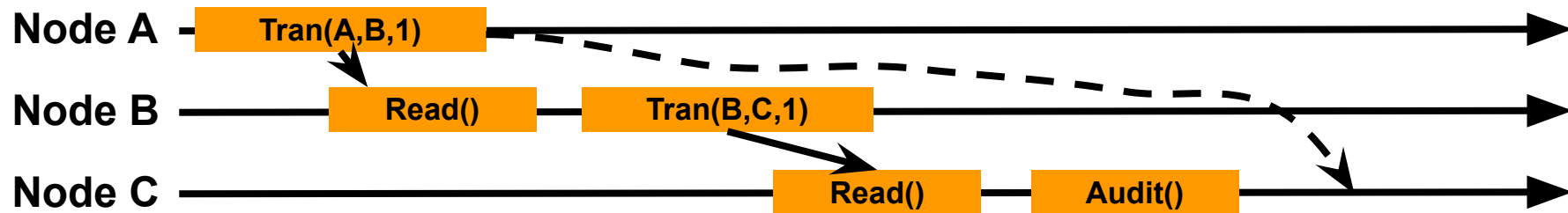
[Ahamad et al. DC 95]

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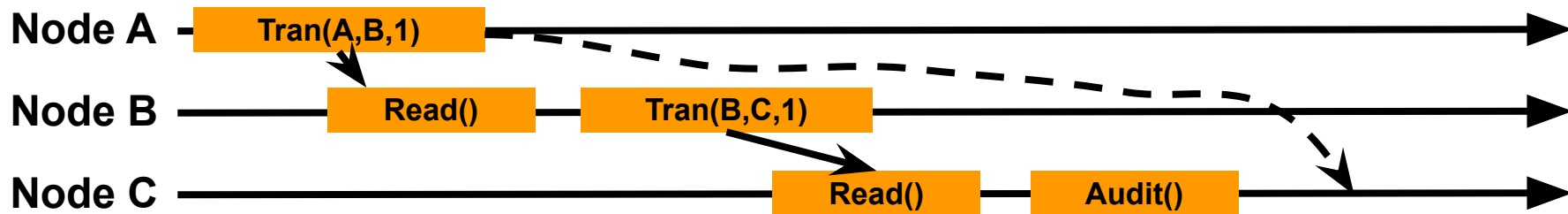
**Write1** → **Read1** → **Write2** → **Read2**  
⇒ **Write1** → **Read2**

## Thm2: Causal consistency is necessary





## Thm2: Causal consistency is necessary



Audit() does NOT contain Tran(A,B,1)

⇒ Node C violates auditability!

# Positive Results

## **CCC: Causal Cryptocurrency under Crash faults**

- Similar to ASO-Crypto, but use causal memory underneath

## **CCB: Causal Cryptocurrency under Byzantine faults**

- Byzantine causal memory [Tseng et al. NCA 19]
- Reliable broadcast [Bracha and Toueg JACM 85]
- Sequence number to stop double-spending
- PKI and digital signature
- A weaker form of eventual delivery:  
one needs to be able to talk to  $n-f$  correct nodes

# Summary

ASO-Crypto: consensus not necessary

Our work: total order and strong consistency not necessary

causal consistency necessary

an inherent challenge of Byzantine crypto in partition

two implementations

# Future Works

Implementation and evaluation

Permission-less systems

Probabilistic guarantees

# Advice

Know fundamentals

- FLP
- CAP

Reach out to other communities

Be comfortable with formalism

# Thanks!

# Questions?

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